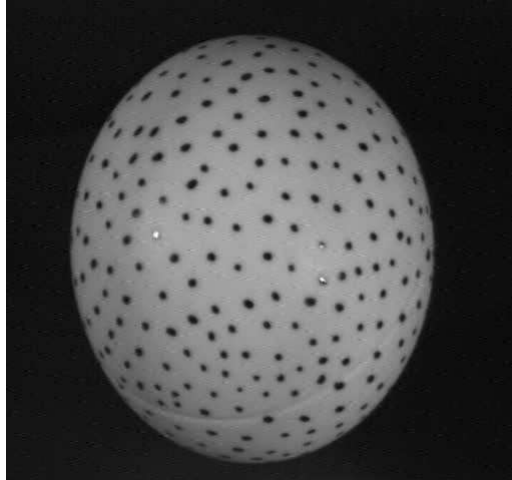


## Motion Tracking

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## Motion tracking

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Suppose we have more than two images

- How to track a point through all of the images?
  - In principle, we could estimate motion between each pair of consecutive frames
  - Given point in first frame, follow arrows to trace out it's path
  - Problem: DRIFT
    - » small errors will tend to grow and grow over time—the point will drift way off course

### Feature Tracking

- Choose only the points (“features”) that are easily tracked
- How to find these features?
  - windows where  $\sum \nabla I (\nabla I)^T$  has two large eigenvalues
- Called the Harris Corner Detector

## Feature Detection

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## Tracking features

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### Feature tracking

- Compute optical flow for that feature for each consecutive H, I

### When will this go wrong?

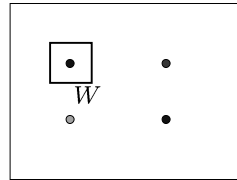
- Occlusions—feature may disappear
  - need mechanism for deleting, adding new features
- Changes in shape, orientation
  - allow the feature to deform
- Changes in color
- Large motions
  - will pyramid techniques work for feature tracking?

## Handling large motions

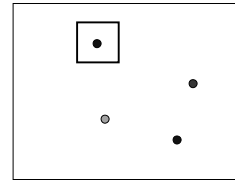
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L-K requires small motion

- If the motion is much more than a pixel, use discrete **search** instead



$H(x, y)$



$I(x, y)$

- Given feature window  $W$  in  $H$ , find best matching window in  $I$
- Minimize sum squared difference (SSD) of pixels in window

$$\min_{(u,v)} \left\{ \sum_{(x,y) \in W} |I(x+u, y+v) - H(x, y)|^2 \right\}$$

- Solve by doing a search over a specified range of  $(u,v)$  values
  - this  $(u,v)$  range defines the **search window**

## Tracking Over Many Frames

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### Feature tracking with $m$ frames

1. Select features in first frame
2. Given feature in frame  $i$ , compute position in  $i+1$
3. Select more features if needed
4.  $i = i + 1$
5. If  $i < m$ , go to step 2

### Issues

- Discrete search vs. Lucas Kanade?
  - depends on expected magnitude of motion
  - discrete search is more flexible
- Compare feature in frame  $i$  to  $i+1$  or frame 1 to  $i+1$ ?
  - affects tendency to drift..
- How big should search window be?
  - too small: lost features. Too large: slow

## Incorporating Dynamics

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### Idea

- Can get better performance if we know something about the way points move
- Most approaches assume constant velocity

$$\dot{\mathbf{x}}_{i+1} = \dot{\mathbf{x}}_i$$

$$\mathbf{x}_{i+1} = 2\mathbf{x}_i - \mathbf{x}_{i-1}$$

or constant acceleration

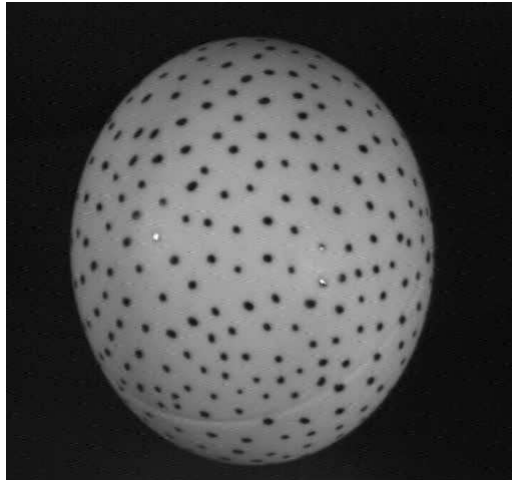
$$\ddot{\mathbf{x}}_{i+1} = \ddot{\mathbf{x}}_i$$

$$\mathbf{x}_{i+1} = 3\mathbf{x}_i - 3\mathbf{x}_{i-1} + \mathbf{x}_{i-2}$$

- Use above to predict position in next frame, initialize search

## Point Tracking (Similar Features)

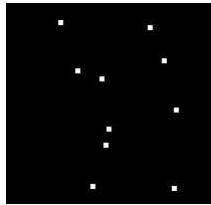
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## Point Tracking

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- All objects are similar
- Only Motion information is available

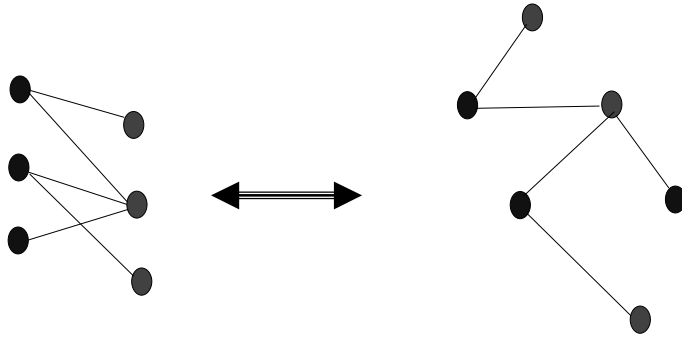


## Graph Theory (Again)

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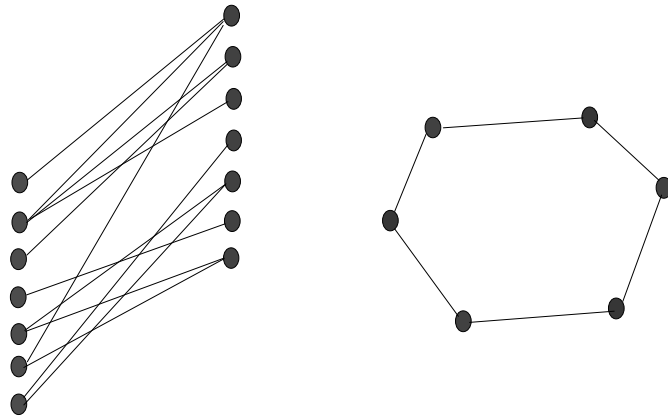
## Graph

A graph  $G(V,E)$  is a triple consisting of a vertex set  $V(G)$  an edge set  $E(G)$  and a relation that associates with each edge two vertices called its end points.



## Bipartite Graph

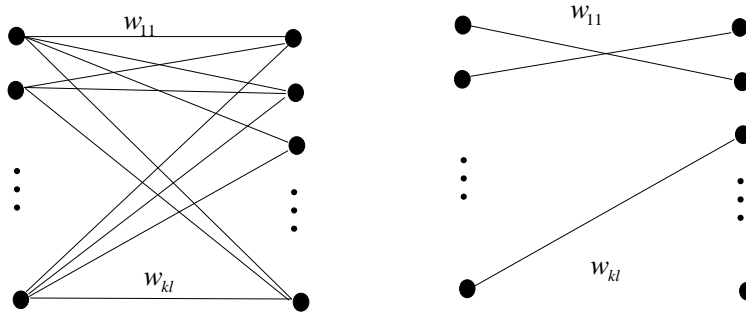
A graph  $G$  is bipartite if its vertex set can be partitioned in two subsets in such a way that no two vertex in same set have a common edge.



## Matching

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Matching is a set of edges such that no two of them have a common vertex



## Point Tracking

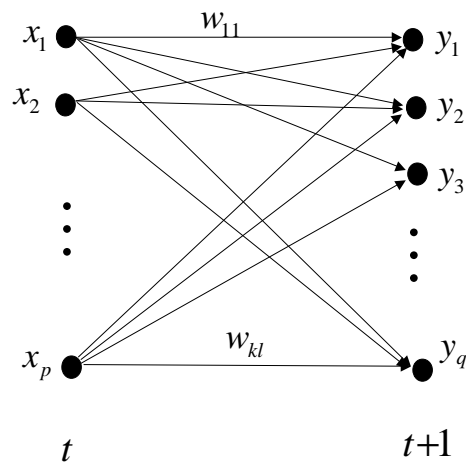
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- Points corresponds to vertices in the bipartite graph
- Points at time instants  $t$  and  $t+1$  form partite sets of graph.
- The cost of corresponding a point at instant  $t$  to a point at instant  $t+1$  is the weight of edge between the corresponding vertices.

S. Ullman, "The interpretation of visual motion," MIT Press, Cambridge, MA.

## Point Tracking

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### Find Minimum Matching of Bipartite Graph

Maximum (Minimum) Matching is a set of edges such that no two of them have a common vertex and the sum of weights is maximum (minimum) among all such sets

## Algorithm

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- Compute costs  $w_{ij}$  for each pair of points
- Construct a bipartite graph based on computed costs
- Prune all edges having weights exceeding certain threshold
- Find the minimum matching of constructed graph. (Hungarian Algorithm)



## Cost Computation

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Ullman:  $w_{ij} = \|y_j - x_i\|$

Sethi & Jain:  $w_{ij} = c \left[ 1 - \frac{(y_j - x_i) \cdot v_{x_i}}{\|y_j - x_i\| \|v_{x_i}\|} \right] + (1 - c) \left[ 1 - \frac{GM(\|y_j - x_i\|, \|v_{x_i}\|)}{AM(\|y_j - x_i\|, \|v_{x_i}\|)} \right]$

$$GM(a, b) = \sqrt{ab}$$

$$AM(a, b) = \frac{a + b}{2}$$

Rangarajan & Shah:  $w_{ij} = \frac{\|v_{x_i} - (y_j - x_i)\|}{\sum_{k=1}^p \sum_{l=1}^q \|v_{x_k} - (y_l - x_k)\|} + \frac{\|y_j - x_i\|}{\sum_{k=1}^p \sum_{l=1}^q \|y_l - x_p\|}$

## Can we increase the search space further in time?

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- An N-Dimensional Matching is a set  $M$  of hyper-edges or s-tuples such that no two hyper-edges of  $M$  have common vertices.
- A minimum N-D Matching is an N-D Matching with minimum weight
- Finding a minimum N-D Matching is NP-Hard Problem

## Change Detection

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C. Stauffer and W.E.L. Grimson, "Learning patterns of activity using real time tracking,"  
IEEE Trans. On PAMI, 22(8):747-757, Aug 2000