

# Lecture-7

## Step Length Selection

### Homework (Due 2/20/01)

- 3.1
- 3.2
- 3.5
- 3.6
- 3.7
- 3.9
- 3.10
- Show equation 3.44
- The last step in the proof of Theorem 3.6. (see slides)

## Sufficient condition

$$f(x_k + \mathbf{a}p_k) \leq f(x_k) + c_1 \mathbf{a} \nabla f_k^T p_k, \quad c_1 \in (0,1) \quad c_1 = 10^{-4}$$

$$f(x_k + \mathbf{a}p_k) - f(x_k) \leq c_1 \mathbf{a} \nabla f_k^T p_k, \quad c_1 \in (0,1)$$

The reduction should be proportional to both the step length, and directional derivative.

$$f(x_k + \mathbf{a}p_k) \leq f(x_k) + c_1 \mathbf{a} \nabla f_k^T p_k, \quad c_1 \in (0,1)$$

$$f(x_k + \mathbf{a}p_k) \leq l(\mathbf{a})$$

St line

## Sufficient condition

$$f(x_k + \mathbf{a}p_k) \leq l(\mathbf{a})$$

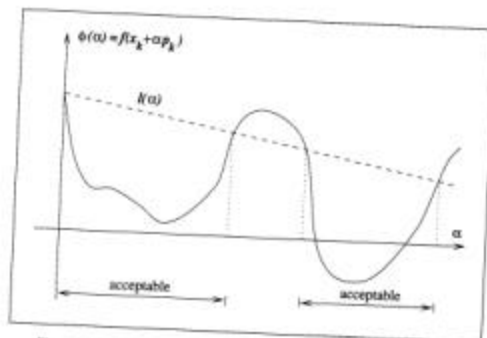


Figure 3.3 Sufficient decrease condition.

Problem:  
The sufficient decrease condition is satisfied for all small values of step length

## Curvature condition

$$\nabla f(x_k + \alpha p_k)^T p_k \geq c_2 \nabla f_k^T(x_k) p_k, \quad c_2 \in (c_1, 1)$$

Derivative

$c_2 = .9$  for Newton and Quasi-Newton

$c_2 = .1$  for conjugate gradient

The slope of  $\nabla f(x_k + \alpha p_k)^T p_k$  is greater than  $c_2$  times the gradient  $\nabla f_k^T(x_k) p_k$ .

## Curvature condition

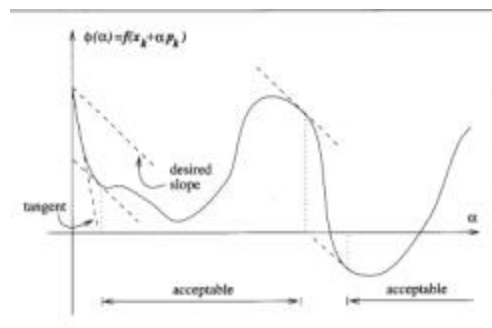


Figure 3.4 The curvature condition.

If the slope is strongly negative, that means we can reduce  $f$  further along the chosen direction

If the slope is positive, it indicates we can not decrease  $f$  further in this direction.

## Wolfe conditions

$$f(x_k + \mathbf{a}p_k) \leq f(x_k) + c_1 \mathbf{a} \nabla f_k^T p_k, \quad c_1 \in (0,1) \quad \text{Sufficient decrease}$$

$$\nabla f(x_k + \mathbf{a}p_k)^T p_k \geq c_2 \nabla f_k^T p_k, \quad c_2 \in (c_1,1) \quad \text{Curvature}$$

## Backtracking Line Search

If line search method chooses its step length appropriately, we can dispense with the second condition

Choose  $\bar{\mathbf{a}} > 0$ ,  $\mathbf{r}, c \in (0,1)$ ; set  $\mathbf{a} \leftarrow \bar{\mathbf{a}}$ ;

*repeat* until  $f(x_k + \mathbf{a}p_k) \leq f(x_k) + c \mathbf{a} \nabla f_k^T p_k$

$\mathbf{a} \leftarrow \mathbf{r} \mathbf{a}$ ;

*end(repeat)*

Terminate with  $\mathbf{a}_k = \mathbf{a}$

$\bar{\mathbf{a}} = 1$ , for Newton  
and quasi - Newton

This ensures that the step length is short enough to satisfy the sufficient decrease condition, but not too short.

## Searching Step Length Using Interpolation

$$f(x_k + \alpha p_k) \leq f(x_k) + c_1 \alpha \nabla f_k^T p_k, \quad c_1 \in (0,1) \quad \text{Sufficient decrease}$$

$$f(x_k) \leq f(0) + c_1 \alpha_k f'(0)$$

1. Assume  $a_0$  is the initial guess. Then if we have:

$$f(a_0) \leq f(0) + c_1 \alpha_0 f'(0)$$

Then this step length satisfies the condition, we terminate the search.

2. Otherwise, we know  $[0, a_0]$  contains the acceptable step lengths.

We fit quadratic polynomial to three pieces of information:

$$f_q(0) = f(0), f'_q(0) = f'(0), f_q(a_0) = f(a_0)$$

## Searching Step Length Using Interpolation

and find step length  $a_1$  by analytically minimizing this polynomial

If the sufficient decrease condition is satisfied for this  $a_1$  then we terminate the search.

If not we fit cubic polynomial to interpolate four pieces of information, and analytically minimize this polynomial to find  $a_1$ .

$$f_c(0) = f(0), f'_c(0) = f'(0), f_c(a_0) = f(a_0), f_c(a_1) = f(a_1)$$

3. If not we fit cubic polynomial to interpolate four pieces of  $a_2$  information, and analytically minimize this polynomial to find  $a_2$ .

If necessary we can repeat this process with  $f(0), f'(0)$  and two Most recent values of  $f$ .

## Quadratic Interpolation

$$f_q(\mathbf{a}) = a\mathbf{a}^2 + b\mathbf{a} + c$$

$$f_q(0) = f(0), f'_q(0) = f'(0), f_q(\mathbf{a}_0) = f(\mathbf{a}_0)$$

$$f_q(\mathbf{a}) = \left( \frac{f(\mathbf{a}_0) - f(0) - \mathbf{a}_0 f'(0)}{\mathbf{a}_0^2} \right) \mathbf{a}^2 + f'(0)\mathbf{a} + f(0)$$

$$\frac{d}{d\mathbf{a}} f_q(\mathbf{a}) = 2 \left( \frac{f(\mathbf{a}_0) - f(0) - \mathbf{a}_0 f'(0)}{\mathbf{a}_0^2} \right) \mathbf{a} + f'(0) = 0$$

$$\mathbf{a} = - \left( \frac{f(0)\mathbf{a}_0^2}{2(f(\mathbf{a}_0) - f(0) - \mathbf{a}_0 f'(0))} \right)$$

## Cubic Interpolation

$$f_c(\mathbf{a}) = a\mathbf{a}^3 + b\mathbf{a}^2 + c\mathbf{a} + d$$

$$f_c(0) = f(0), f'_c(0) = f'(0), f_c(\mathbf{a}_0) = f(\mathbf{a}_0), f'_c(\mathbf{a}_0) = f'(\mathbf{a}_0)$$

$$f_c(\mathbf{a}) = a\mathbf{a}^3 + b\mathbf{a}^2 + f'(0)\mathbf{a} + f(0)$$

$$\begin{bmatrix} a \\ b \end{bmatrix} = \frac{1}{\mathbf{a}_0^2 \mathbf{a}_1^2 (\mathbf{a}_1 - \mathbf{a}_0)} \begin{bmatrix} \mathbf{a}_1^2 & -\mathbf{a}_0^2 \\ -\mathbf{a}_0^3 & \mathbf{a}_1^3 \end{bmatrix} \begin{bmatrix} f(\mathbf{a}_1) - f(0) - f'(0)\mathbf{a}_1 \\ f(\mathbf{a}_0) - f(0) - f'(0)\mathbf{a}_0 \end{bmatrix}$$

$$\mathbf{a}_2 = - \left( \frac{-b + \sqrt{b^2 - 3af'(0)}}{3a} \right)$$

## Algorithm 3.2 (Line Search Algorithm)

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Set  $\mathbf{a}_0 \leftarrow 0$ , choose  $\mathbf{a}_1 > 0$ , and  $\mathbf{a}_{\max}$ ;
 $i \leftarrow 1$ 
repeat
  Evaluate  $f(\mathbf{a})$ ;
  if  $f(\mathbf{a}) > f(0) + c_1 \mathbf{a} f'(0)$  or  $[f(\mathbf{a}) > f(\mathbf{a}_{i-1}), i > 1]$  1st Wolfe's condition
     $\mathbf{a} \leftarrow \text{zoom}(\mathbf{a}_{i-1}, \mathbf{a})$ , and stop;
  Evaluate  $f(\mathbf{a})$ ;
  if  $|f'(\mathbf{a})| \leq -c_2 f'(0)$  2nd Wolfe's condition
    set  $\mathbf{a} \leftarrow \mathbf{a}$ , and stop;
  if  $f(\mathbf{a}) \geq 0$ 
    set  $\mathbf{a} \leftarrow \text{zoom}(\mathbf{a}, \mathbf{a}_{i-1})$ , and stop;
  choose  $\mathbf{a}_{i+1} \in (\mathbf{a}, \mathbf{a}_{\max})$ 
   $i \leftarrow i + 1$ ;
end(repeat)

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## Algorithm 3.3 (Zoom)

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repeat
  Interpolate to find a trial step length
   $\mathbf{a}_j$  between  $\mathbf{a}_{lo}, \mathbf{a}_{hi}$ ;
  Evaluate  $f(\mathbf{a}_j)$ ;
  if  $f(\mathbf{a}_j) > f(0) + c_1 \mathbf{a}_j f'(0)$  or  $[f(\mathbf{a}_j) > f(\mathbf{a}_{lo})]$  1st Wolfe's condition
     $\mathbf{a}_{hi} \leftarrow \mathbf{a}_j$ ;
  else
    Evaluate  $f(\mathbf{a}_j)$ ;
    if  $|f'(\mathbf{a}_j)| \leq -c_2 f'(0)$  2nd Wolfe's condition
      set  $\mathbf{a}_* \leftarrow \mathbf{a}_j$ , and stop;
    if  $f'(\mathbf{a}_j)(\mathbf{a}_{hi} - \mathbf{a}_{lo}) \geq 0$ 
      set  $\mathbf{a}_{hi} \leftarrow \mathbf{a}_{lo}$ 
    set  $\mathbf{a}_{lo} \leftarrow \mathbf{a}_j$ 
  end(repeat)

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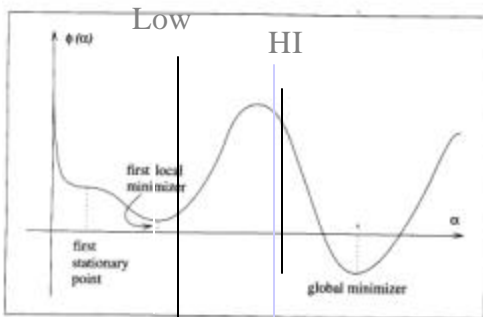


Figure 3.1 The ideal step length is the global minimizer

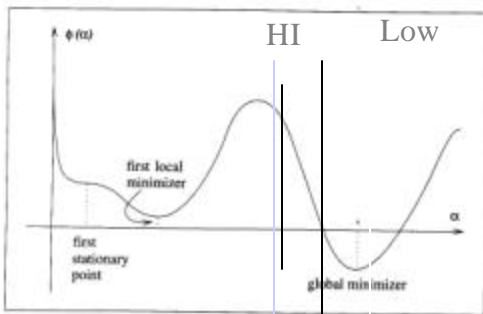


Figure 3.1 The ideal step length is the global minimizer



## Theorem 3.5 (Any Descent Direction)

Suppose  $f$  is three times continuously differentiable. Consider iteration  $x_{k+1} = x_k + \alpha_k p_k$ , where  $p_k$  is a descent direction, Satisfies Wolfe's conditions, with  $\alpha_k > 0$ . If the sequence  $\{x_k\}$  converges to a point  $x^*$  such that  $\nabla^2 f(x^*)$  is pd, and if the search direction satisfies

$$\lim_{k \rightarrow \infty} \frac{\|\nabla f_k + \nabla^2 f_k p_k\|}{\|p_k\|} = 0$$

$$\lim_{k \rightarrow \infty} \frac{\|(B_k - \nabla^2 f(x^*))p_k\|}{\|p_k\|} = 0$$

Then

- (i)  $\alpha_k$  is admissible for all  $k > k_0$  and
- (ii) if  $\alpha_k \rightarrow 0$  for all  $k > k_0$ , then  $\{x_k\}$  converges to  $x^*$  superlinearly.

## Theorem 3.6 (Quasi-Newton)

Suppose  $f$  is three times continuously differentiable. Consider iteration  $x_{k+1} = x_k + \alpha_k p_k$ , where  $p_k$  is given by Quasi-Newton direction. Assume the sequence  $\{x_k\}$  converges to a point  $x^*$  such that  $\nabla^2 f(x^*)$  is pd, the sequence  $\{x_k\}$  converges superlinearly iff the following condition holds.

$$\lim_{k \rightarrow \infty} \frac{\|(B_k - \nabla^2 f(x^*))p_k\|}{\|p_k\|} = 0$$

## Order Notations

Given two non-negative infinite sequences

$$\mathbf{h}_k = O(\mathbf{n}_k)$$

$$\text{if } |\mathbf{h}_k| \leq C |\mathbf{n}_k|, \text{ for } C > 0, \forall k$$

$$\mathbf{h}_k = o(\mathbf{n}_k)$$

$$\text{if } \lim_{k \rightarrow \infty} \frac{\mathbf{h}_k}{\mathbf{n}_k} = 0$$

## Sketch of a Proof

$$\begin{aligned} p_k - p_k^N &= \nabla^2 f_k^{-1} (\nabla^2 f_k p_k + \nabla f_k) \\ &= \nabla^2 f_k^{-1} (\nabla^2 f_k - B_k) p_k \\ &= O(\| (\nabla^2 f_k - B_k) p_k \|) \\ &= o(\| p_k \|) \end{aligned}$$

$$\mathbf{h}_k = O(\mathbf{n}_k)$$

$$\text{if } |\mathbf{h}_k| \leq C |\mathbf{n}_k|, \text{ for } C > 0$$

$$\mathbf{h}_k = o(\mathbf{n}_k)$$

$$\text{if } \lim_{k \rightarrow \infty} \frac{\mathbf{h}_k}{\mathbf{n}_k} = 0$$

$$\lim_{k \rightarrow 0} \frac{\| (B_k - \nabla^2 f(x^*)) p_k \|}{\| p_k \|} = 0$$

$$\lim_{k \rightarrow 0} \frac{\| \nabla f_k + \nabla^2 f_k p_k \|}{\| p_k \|} = 0$$

Norm of Hessian is bounded.

## Sketch of a Proof

$$\begin{aligned} \|x_k + p_k - x^*\| &= \|x_k + p_k^N - p_k^N + p_k - x^*\| \leq \|x_k + p_k^N - x^*\| + \|p_k - p_k^N\| \\ &= O(\|x_k - x^*\|^2) + o(\|p_k\|) \\ \|x_k + p_k - x^*\| &\leq o(\|x_k - x^*\|) \end{aligned}$$

$$\begin{aligned} h_k &= o(n_k) \\ \text{if } \lim_{k \rightarrow \infty} \frac{h_k}{n_k} &= 0 \end{aligned}$$

Theorem 3.7

Super-linear

Show this in Homework

## Theorem 3.7 (Newton)

Suppose that  $f$  is twice differentiable and that Hessian is Lipschitz continuous. Consider the iteration  $p_k^N = -\nabla^2 f_k^{-1} \nabla f_k$  where  $p_k$  is given by

$$p_k^N = -\nabla^2 f_k^{-1} \nabla f_k$$

Then:

1. If the starting point  $x_0$  is sufficiently close to  $x^*$ , the sequence converges to  $x^*$ .
2. The rate of convergence is quadratic
3. The sequence of gradient norms  $\|\nabla f_k\|$  converges quadratically to zero.

## Coordinate Descent Method

Cycle through  $n$  coordinate directions  $e_1, e_2, \dots, e_n$  using each in turn as a search direction.

Fix all other variables except one, and minimize the function.

It is an inefficient method, it can iterate infinitely without ever approaching a point, where the gradient vanishes.

The gradient may become more and more perpendicular to search directions, making  $\cos \mathbf{q}$  approach to zero, but not the gradient.

## Solution of A linear System

- Gaussian Elimination, Backward Substitution
- Matrix Factorization
- Iterative Techniques

**6.1 Gaussian Elimination with Backward Substitution**

To solve the  $n \times n$  linear system

$$\begin{aligned} E_1: & a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = a_{1n+1} \\ E_2: & a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = a_{2n+1} \\ & \vdots \\ E_n: & a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = a_{nn+1} \end{aligned}$$

INPUT: number of unknowns and equations  $n$ ; augmented matrix  $A = (a_{ij})$ ,  $i \leq n$  and  $1 \leq j \leq n+1$ .

OUTPUT: solution  $x_1, x_2, \dots, x_n$ , or message that the linear system has no solution.

Step 1: For  $i = 1, \dots, n-1$  do Steps 2-6. (Elimination process.)

Step 2: Let  $p$  be the smallest integer with  $i \leq p \leq n$  and  $a_{ip} \neq 0$ .  
If no integer  $p$  can be found then OUTPUT ('no unique solution exists'); STOP.

Step 3: If  $p \neq i$  then perform  $(E_p) \leftrightarrow (E_i)$ .

Step 4: For  $j = i+1, \dots, n$  do Steps 5 and 6.

Step 5: Set  $m_{ij} = a_{ij}/a_{ip}$ .

Step 6: Perform  $(E_i - m_{ij}E_j) \rightarrow (E_i)$ .

Step 7: If  $a_{nn} = 0$  then OUTPUT ('no unique solution exists'); STOP.

Step 8: Set  $x_n = a_{nn+1}/a_{nn}$ . (Start backward substitution.)

Step 9: For  $i = n-1, \dots, 1$  set  $x_i = [a_{i,n+1} - \sum_{j=i+1}^n a_{ij}x_j] / a_{ii}$ .

Step 10: OUTPUT  $(x_1, \dots, x_n)$ . (Procedure completed successfully.) STOP.

## Iterative Methods for Solving Linear Systems

- For large sparse system Gaussian Elimination and Backward substitution is not suitable.
- Approximate solution using iterative methods

## Jacobi

$$x_i^k = \frac{\sum_{j=1, j \neq i}^n (-a_{ij} x_j^{k-1}) + b_i}{a_{ii}} \quad \text{for } i = 1, 2, \dots, n$$

## Gauss-Seidel

$$x_i^k = \frac{-\sum_{j=1}^{i-1} (a_{ij} x_j^k) - \sum_{j=i+1}^n (a_{ij} x_j^{k-1}) + b_i}{a_{ii}} \quad \text{for } i = 1, 2, \dots, n$$

## SOR (Successive Over Relaxation)

$$x_i^k = (1 - \mathbf{w})x_i^{k-1} + \frac{\mathbf{w}}{a_{ii}} \left[ b_i - \sum_{j=1}^{i-1} a_{ij}x_j^k - \sum_{j=i+1}^n (a_{ij}x_j^{k-1}) \right]$$

$$\mathbf{w} > 1$$