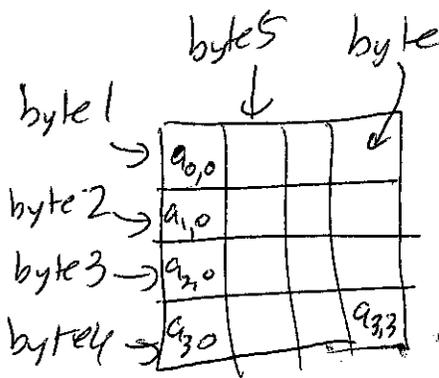


Rijndael (pronounced RAIN-DOLL)

- SECURE
 - FAST
 - SIMPLE
 - EXTENDIBLE (128 bit key, 192 bit key, 256 bit key)
- operations are very fast in either software or hardware



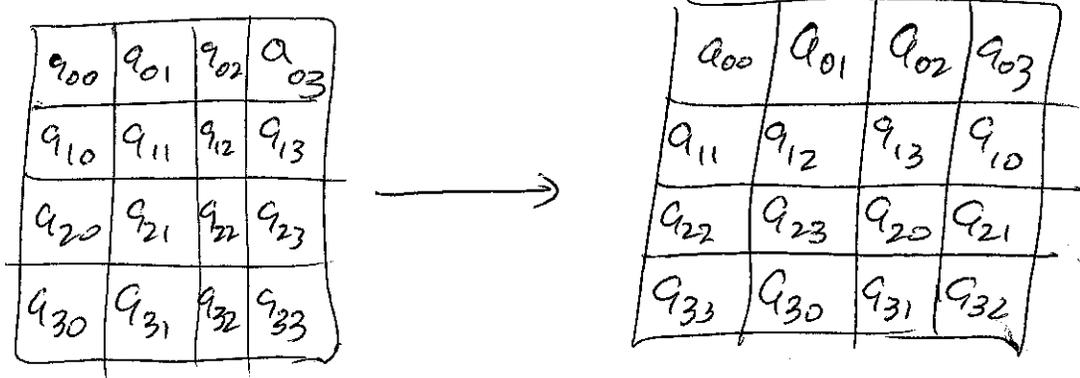
16 squares x 8 bits = 128 bits
BLOCK SIZE

Generally bytes expressed in HEX

7C = 01111100

S-box(7C) = 10, S-inv(10) = 7C
S-box(A4) = 49, S-inv(49) = A4

Input (SHIFT ROWS)



Mix Columns

$$\text{row 2} \begin{bmatrix} 2 & 3 & 1 & 1 \\ 1 & 2 & 3 & 1 \\ 1 & 1 & 2 & 3 \\ 3 & 1 & 1 & 2 \end{bmatrix} \begin{bmatrix} 74 & 8F & 26 & 39 \\ CB & 04 & D1 & EE \\ 67 & 4B & 39 & 48 \\ BF & F1 & 00 & 32 \end{bmatrix} = \begin{bmatrix} - & - & - & - \\ - & - & - & - \\ - & - & - & - \\ - & - & - & - \end{bmatrix}$$

col 4

each byte is actually a special type of polynomial.

$$39 = 00111001 = x^5 + x^4 + x^3 + 1$$

→ poly is deg 7 or less, coefficients 0,1
polynomials for AES for a field in $GF(2^8)$.

all calculations are done mod $x^8 + x^4 + x^3 + x + 1$.

if $f(x)$ is degree 7 or less then

$$f(x) \text{ mod } x^8 + x^4 + x^3 + x + 1 = f(x).$$

Irreducible
Polynomial

Since only coefficients are 0 and 1,
all coefficients are considered mod 2.

$$\begin{aligned} & (x^6 + x^5 + x^3 + x + 1) + (x^7 + x^5 + x^3 + x) \\ & = \boxed{x^7 + x^6 + 1} \end{aligned}$$

$$(x^4 + 1)(x^4 + x^2 + x)$$

$$= x^8 + x^6 + x^5 + x^4 + x^2 + x$$

$$\begin{array}{r}
 x^8 + x^4 + x^3 + x + 1 \\
 - (x^8 + x^6 + x^5 + x^4 + x^2 + x) \\
 \hline
 x^6 + x^5 + x^3 + x^2 + 1
 \end{array}$$

If I have 1 in the 9th bit position, the 1 is equivalent to getting rid of the 1 and flipping bits $x^4, x^3, x, 1$.

$$\begin{array}{r}
 1011101\text{~~01~~} \\
 \oplus \quad 11011 \\
 \hline
 011011\text{~~01~~} \\
 \quad \quad \quad 01
 \end{array}$$

2	3	1	1
1	2	3	1
1	1	2	3
3	1	1	2

7A	8F	26	39
CB	04	D1	EE
67	4B	39	48
BF	F1	00	32

$1 \times 39 = 39$	$= 00111001$
$2 \times EE = C7$	$= 11000111$
$3 \times 48 = D8$	$= 11011000$
$1 \times 32 = 32$	$= 00110010$
	$\hline 00010100 = 14$

mult by 2 is just adding 0.

$$\begin{array}{r}
 (10) \times (11101110) \\
 \times \\
 \hline
 111011100 \\
 \oplus 11011100 \\
 \hline
 11000111
 \end{array}$$

$$\hookrightarrow (11) \times (01001000) =$$

\downarrow
x+1

$$\begin{array}{r} 01001000 \\ \oplus 10010000 \\ \hline \end{array}$$

$$11011000$$

10/3/16

~~3~~
4

Most of the time

10/5/16 ①

$$W[i] = W[i-4] \oplus W[i-1]$$

If triggers when we start to fill in a new round key. → Uses S-box

$$\text{temp} = \text{subword}(\text{rotword}(\text{temp})) \oplus \text{Rcon}[i/4]$$

↓
left cyclic rotation by one byte

↓
pad with 3 bytes of 0s on right.

$$W[35] = 36A7490F$$

$$\text{temp} = W[35] = AB286C5D$$

$$W[36] = ?$$

$$\text{rotword}(\text{temp}) = 286C5DAB$$

$$\text{subword}(\downarrow) = 34504C62$$

$$\text{Rcon}[i/4] \oplus 1B000000$$

$$\hline 2F504C62$$

$$\oplus 36A7490F$$

$$\hline W[36] = 19F7056D$$

$$\begin{array}{r} 100 \\ 1011 \\ \hline 1111 \end{array}$$

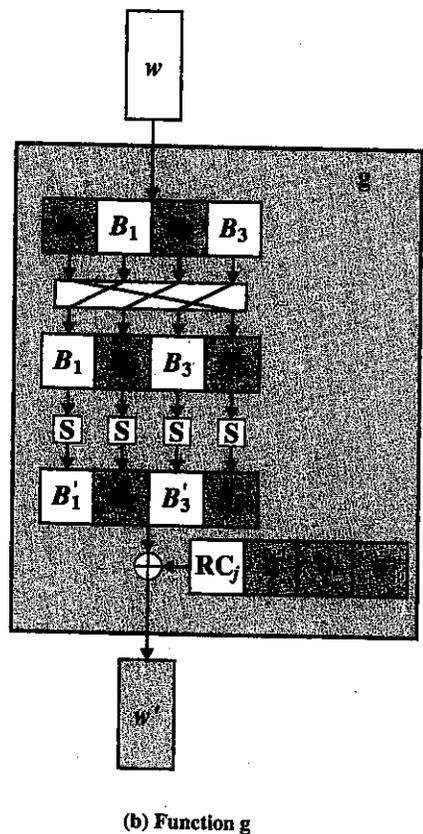
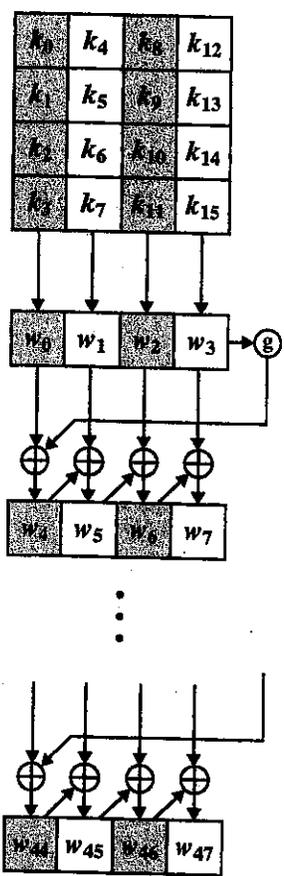
$$\begin{array}{r} 1111 \\ 0110 \\ 1010 \\ 0101 \\ \hline 1100 \\ 1001 \\ \hline 0101 \\ 1111 \\ 10 \end{array}$$

```

KeyExpansion (byte key[16], word w[44])

word temp;
for (i = 0; i < 4; i++) w[i] = (key[4*i], key[4*i+1],
                                   key[4*i+2], key[4*i+3]);

for (i = 4; i < 44; i++)
{
    temp = w[i-1];
    if (i mod 4 = 0) temp = SubWord (RotWord (temp))
                      ⊕ RCn[i/4];
    w[i] = w[i-4] ⊕ temp;
}
    
```



(a) Overall algorithm

(b) Function g

Figure 5.9 AES Key Expansion

tion defined over the field $GF(2^8)$. The values of $RC[j]$ in hexadecimal are

j	1	2	3	4	5	6	7	8	9	10
RC[j]	01	02	04	08	10	20	40	80	1B	36