

8/28/2020 - Extended Euclidean Algorithm

Divisibility and Mod Definitions

We say that $a \mid b$, "b is divisible by a" if and only if, there exists some integer c such that $b = ac$.

So true statements are $6 \mid 18$, $22 \mid 22$, $173 \mid 0$, $7 \mid 98$, etc.

When I say that $a \equiv b \pmod{n}$, this means that $n \mid (a - b)$. (Intuitively, it's like saying that a and b yield the same remainder when divided by n.) Some true mod statements are:

$13 \equiv 5 \pmod{8}$	because $8 \mid (13 - 5)$
$127 \equiv 7 \pmod{10}$	because $10 \mid (127 - 7)$
$127 \equiv 107 \pmod{10}$	because $10 \mid (127 - 107)$
$-5 \equiv 21 \pmod{26}$	because $26 \mid (-5 - 21)$
$-5 \equiv 73 \pmod{26}$	because $26 \mid (-5 - 73)$

In general if we know that $a \equiv b \pmod{n}$, a is also equivalent to $(b + \text{any multiple of } n) \pmod{n}$. so, we can get different equivalent mods by adding or subtracting multiples of n.

$21 \pmod{26}$, some equivalent mods are -5, 47, 73, 99, etc.

Affine Cipher

Encryption Function $f(x) = (ax + b) \pmod{26}$, keys are a, and b.

a has to be from the set $\{1,3,5,7,9,11,15,17,19,21,23,25\}$

b is any int from 0 to 25.

312 possible keys.

Why are there only some possible values of a?

$$F(x) = (ax + b) \pmod{26}$$

So, to do this backwards, we must solve for x in this equation. So usually when we invert functions, we swap x and y and then solve for y:

$$x \equiv (ay + b) \pmod{26}$$

$$(x - b) \equiv ay \pmod{26}$$

Now we are stuck because we are NOT allowed to divide under mod.

But for some values of a, we can multiply the equation through and solve for y anyway...

$$(x - 19) \equiv 5y \pmod{26} \text{ (for example } a = 5, b = 19)$$

$$21(x - 19) \equiv 21(5y) \pmod{26}$$

$$21x - 399 \equiv 105y \pmod{26}$$

Since $105 \equiv 1 \pmod{26}$, then we can say

$$21x - 399 \equiv y \pmod{26}$$

$$y = (21x - 399) \pmod{26}$$

$$y = (21x + 17) \pmod{26}$$

Question: How do I find the magic value? Also, when does the magic value exist and when doesn't it exist?

General form of what we are trying to solve is

$$ay \equiv (x - b) \pmod{n}$$

Here I use n instead of 26 to allow for different alphabet sizes.

What we are really looking for is some value a', such that

$$aa' \equiv 1 \pmod{n}$$

Consider a situation where a and n have some common factor c, where $c > 1$.

This means there is some integer d such that $a = dc$, and there is some integer e such that $n = ec$

$$(dc)a' \equiv 1 \pmod{ec}$$

Let's consider the difference between dca' and 1.

$dca' - 1$, can NOT have a factor of c , because c is greater than 1, and we can pull it out of the first term but not the second. This isn't divisible by c .

But, ec IS divisible by c .

If two things are going to be equivalent mod ec , they must be equivalent mod c , but they aren't which means that the equation is impossible to satisfy.

$$4x \equiv 1 \pmod{26}$$

There is no value of x that satisfies this equation...the reason is that the LHS is always even and the right hand side is always odd...

$$30x \equiv 1 \pmod{999}$$

In this situation the LHS is always divisible by 3, but the RHS is never divisible by 3.

if $ax \equiv 1 \pmod{n}$, then if a and n share a common factor, say c ,

Then the lefthand side is divisible by c , but the right hand side is NOT divisible by c , which means there is no possible integer value of x which makes the equation true.

What we've proven so far, AND THIS IS CRITICAL, is that we CAN NOT INVERT the affine cipher if there is a common factor greater than 1 between a and n , where n is the alphabet size.

**Step 2: Can we find that magic value if a and n share no common factors?
If so, how?**

Answer: Yes, the how is by using the Extended Euclidean Algorithm

Regular Euclidean Algorithm

Given positive integer values, a and b, find the greatest common divisor of a and b.

a = 105, b = 39, idea is to divide a by b, and get the quotient and remainder.

$$105 = 2 \times 39 + 27$$

$$39 = 1 \times 27 + 12$$

$$27 = 2 \times 12 + 3, \text{ this is the gcd}(105, 39) = 3, \text{ the last non-zero remainder in the}$$

$$12 = 4 \times 3 \quad \text{process.}$$

This means, for example, that there is no solution to

$$39x \equiv 1 \pmod{105}$$

Let's say we have values of a and b that don't share a common factor:

$$a = 143, b = 105$$

$$143 = 1 \times 105 + 38$$

$$143 - 1 \times 105 = 38$$

$$105 = 2 \times 38 + 29$$

$$105 - 2 \times 38 = 29$$

$$38 = 1 \times 29 + 9$$

$$38 - 1 \times 29 = 9$$

$$29 = 3 \times 9 + 2$$

$$29 - 3 \times 9 = 2$$

$$9 = 4 \times 2 + 1, \text{ greatest common divisor is 1.}$$

$$2 = 2 \times 1$$

Now, we are interested in finding the answer to

$$105x \equiv 1 \pmod{143}$$

$$\underline{9} - 4 \times \underline{2} = 1$$

$$9 - 4(\underline{29} - 3 \times \underline{9}) = 1$$

$$\underline{9} - 4 \times 29 + \underline{12} \times \underline{9} = 1$$

$$13 \times \underline{9} - 4 \times 29 = 1$$

$$13(\underline{38} - 1 \times \underline{29}) - 4 \times 29 = 1$$

$$13 \times 38 - 13 \times 29 - 4 \times 29 = 1$$

$$13 \times 38 - 17 \times \underline{29} = 1$$

$$13 \times 38 - 17(\underline{105} - 2 \times \underline{38}) = 1$$

$$\underline{13} \times \underline{38} - 17 \times 105 + \underline{34} \times \underline{38} = 1$$

$$47 \times \underline{38} - 17 \times 105 = 1$$

$$\begin{aligned}47(143 - 1 \times 105) - 17 \times 105 &= 1 \\47 \times 143 - 47 \times 105 - 17 \times 105 &= 1 \\47 \times 143 - 64 \times 105 &= 1\end{aligned}$$

Now, let's take this equation mod 143:

$$\begin{aligned}47 \times 143 - 64 \times 105 &\equiv 1 \pmod{143} \\-64 \times 105 &\equiv 1 \pmod{143}\end{aligned}$$

So, our magic value that we multiply 105 by to obtain $1 \pmod{143}$ is -64. Normally, we are going to want to express this as a number in between 1 and 142, so let's add 143 to it to get our final answer: $-64 + 143 = 79$.

$$79 \times 105 \equiv 1 \pmod{143}$$

Goal of the Extended Euclidean Algorithm is to find the value a' such that

$$aa' \equiv 1 \pmod{n}$$

We can do this so long as $\gcd(a, n) = 1$, using the process outlined above.

What we say is that

$$105^{-1} \equiv 79 \pmod{143}$$