

# "Runtime" Comparisons

Which of these grows the fastest

$$\log_2(N^5) \quad N^1 \quad \log(N^n)$$

$$\begin{array}{cccc} \log & 5\log_2(N) & \lfloor N \rfloor & N\log(N) \\ \text{slowest growth} & & & \text{fastest growth} \end{array}$$

Note  $\log(N)$  is faster growing than 1

## Exponential vs Polynomial

$$\begin{array}{ccc} \text{Exponential} & & \text{Polynomial} \\ c^N & & N^k \\ > 1 & & > 0 \end{array}$$

Exponential always grows faster than polynomial

- Derivative of  $c^N$  is  $a c^N$  for some positive  $a$ . We don't care about constant factors and  $a$ .

We don't care about constant factors and  $a$ . We don't care about constant factors and  $a$ .

Derivative of  $N^k$  is  $b N^{k-1}$  for integer  $k$

We will eventually reach  $N^0$

Calculus Def. of  $\mathcal{O}(f(n))$

$f(n) = \mathcal{O}(g(n)) \Leftrightarrow$  for some  $c \in \mathbb{R}^+$   
 $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \leq c$

---

Common Error      Best Case

Don't say

"If we let  $n$  be 1, then the

best case is  $\mathcal{O}(1)$ ."

↙ Bad

def. say so  $n$  tends towards  $\infty$

$n$  should be large!

```

int func7(int a[], int sizea, int b[], int sizeb) {
    int i, j;
    for (i = 0; i < sizea; i++)
        for (j = 0; j < sizeb; j++)
            if (a[i] == b[j])
                return 1;
    return 0;
}

```

Best Worst

$O(1)$   $O(|A| \cdot |B|)$

$O(\frac{|sizea|}{sizeb})$

```

int func8(int a[], int sizea, int b[], int sizeb) {
    int i, j;
    for (i = 0; i < sizea; i++)
        if (binSearch(b, sizeb, a[i]))
            return 1;
    return 0;
}

```

Best Worst

$O(1)$   $O(|A| \log(|B|))$

$\underline{O(n^3)}$

$4n^3 + 4n^2$

$n^3 + n^2$

$(8n^2 + 8n)^{\frac{1}{2}}$

$\underline{int func9(int n)}$

$\underline{int x = 0, i, j;}$

$\underline{for (i = 1; i \leq n * (8 * n + 8); i++) {}$

$\underline{\quad for (j = n / 2; j \leq n; j++) {}$

$\underline{\quad \quad x = x + (n - j);}$

$\underline{\quad }}$

$\underline{}}$

$\underline{return x;}$

$\underline{for(i=0; i < |A|; i++)}$

$\underline{\quad sum += j}$

$\underline{if(binSearch(b, sizeb, a[i]))}$

$\underline{\quad return sum}$

$\underline{}}$

$\underline{return 0;}$

$\underline{O(|A| + \log(|B|))}$

$\underline{\neq O(|A|)}$

$i = 1, 2$

$j = \frac{n}{2}, \frac{n}{2}+1, \frac{n}{2}+2, \dots, n$

$n \text{ or } \frac{n}{2}$

$\frac{n}{2} \text{ total}$

```

int func4(int ** array, int n) {
    int i = 0, j = 0;
    while (i < n) {
        while (j < n && array[i][j] == 1)
            j++;
        i++;
    }
    return j;
}

```

$$N = 30$$

( seconds

$$N = 60$$

25 seconds

$\hookrightarrow O(N)$

```

int func5(int ** array, int n) {
    int i = 0, j;
    while (i < n) {
        j = 0;
        while (j < n && array[i][j] == 1)
            j++;
        i++;
    }
    return j;
}

```

for (int i = 0; i < n; i++)  
 for (int j = i; j < n; j++)  
 sum += j

$O(N \cdot \log(N))$   
 Harmonic Series

$$N \cdot \left( \frac{1}{1} + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right) \approx N \cdot \log(n)$$

```

int func6(int array[], int n) {
    int i, j, sum = 0;
    for (i = 0; i < n; i++)
        for (j = i+1; j < n; j++)
            if (array[i] > array[j])
                sum++;
    return sum;
}

```

i=0 j=[1, 2, ..., n-1]  
 i=1 j=[2, 3, ..., n-1]  
 i=2 j=[3, ..., n-1]

$\sum_{i=0}^{n-1} \sum_{j=i+1}^n$   $\Rightarrow O(n^2)$

$\sum_{i=1}^n \sum_{j=1}^i$   $\Rightarrow O(n^2)$

$\sum_{i=1}^n i = 1 + 2 + 3 + \dots + n-3 + n-2 + n-1$