

COP 3502 10/4/23

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- (1) Put <sup>up</sup> Study Group Guidelines
- (2) Test out run time prediction
- (3) Examples to judging run-times of code segments
- (4) Start w/ recurrence relations

$$O(n^2) \quad n = 50,000 \quad t = 6$$

$$T(n) = cn^2$$

100,000

$$T(50000) = c(50000)^2 = 6 \text{ sec}$$

150,000

$$c = \frac{6 \text{ sec}}{50000^2}$$

$$T(100,000) = \frac{6 \text{ sec}}{50000^2} \times (100000)^2$$

$$= 6 \text{ sec} \times \left(\frac{100000}{50000}\right)^2 = 24 \text{ sec} \quad \left. \begin{array}{l} 27 \\ 21 \end{array} \right\}$$

$$T(150,000) = \frac{6 \text{ sec}}{50000^2} \times (150,000)^2 = 6 \text{ sec} \times \left(\frac{150000}{50000}\right)^2 = 54 \text{ sec}$$

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$$O(n\sqrt{n}) \quad n = 2 \times 10^6 \quad t = 6$$

$$T(n) = cn\sqrt{n}$$

$$T(2 \times 10^6) = c \cdot 2 \times 10^6 \sqrt{2 \times 10^6} = 6 \text{ sec}$$

$$c = \frac{6 \text{ sec}}{2\sqrt{2} \times 10^9}$$

$$\begin{aligned}
 \# n &= 5 \times 10^6 \\
 T(5 \times 10^6) &= \frac{6 \text{ sec}}{2\sqrt{2} \times 10^9} \times \overbrace{5 \times 10^6}^n \times \overbrace{\sqrt{5} \times 10^3}^{\sqrt{n}} \\
 &= \frac{6.5\sqrt{5}}{2\sqrt{2}} = \sim 22 \text{ sec}
 \end{aligned}$$

### More Code Segment Analysis

```

int x = 1
while (x < n)
    x = 2 * x;

```

```

int x = n
while (x > 0)
    printf("%d", x % 2);
    x = x / 2;

```

```

int low = 0, high = n - 1;
while (low < high) {
    int mid = (low + high) / 2;
    if ( )
        low = mid + 1;
    else if ( )
        high = mid - 1;
    else
        return;
}

```

$$\begin{aligned}
 2^k &= n \\
 k &= \log_2 n \quad O(\lg n)
 \end{aligned}$$

$$\log_7 n = \frac{\log_2 n}{\log_2 7} = C \log_2 n$$

$\log_2 7$  is circled. Below it,  $\text{const } O(\log_{b_1} n)$  and  $O(\log_{b_2} n)$  are written, with a bracket labeled "Same" pointing to both.

It's accurate to say an algorithm runs in  $O(3n^2 - 7n)$  time, but it's unconventional because  $O(n^2)$  describes the same thing.