

Backtracking

4/22/22

- Optimization to brute force

Brute Force Examples

(1) Permutations

(2) Combination

(3) Or another arrangement of all possibilities

- list each option

- evaluate each option

Example NOT all perms possible

Visit each location once,
but not all pairs of ~~roads~~ locations
are connected

$n=10$ no road $3 \rightarrow 5$

3, 5, -, -, -, -, -, -, -, - } doomed to fail

↓ 8! wasted work

instead of building all permutations,
skip prefixes that "doomed to fail".

In code

```
for (int i = 0; i < n; i++) {
```

```
    if (used[i]) continue;
```

```
    if (placing i in loc k is doomed to fail) continue;
```

←
BACKTRACKING

```
    used[i] = 1;
```

```
    perm[k] = i;
```

```
    // recursion
```

```
    used[i] = 0;
```

```
}
```

Build subset of toppings via bitwise ops

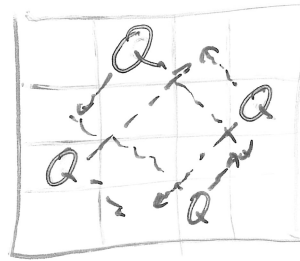
= 00000101 (topping 0, 2)

1 00001000] add topping 3)

→ Storage!

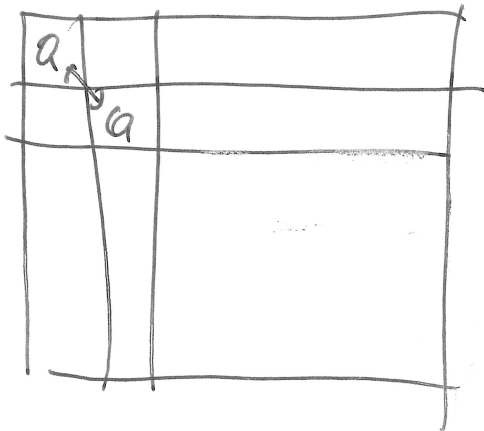
8 Queens

On a chessboard place 8 queens so no two attack each other.



row 0 col 1
row 1 col 3
row 2 col 0
row 3 col 2

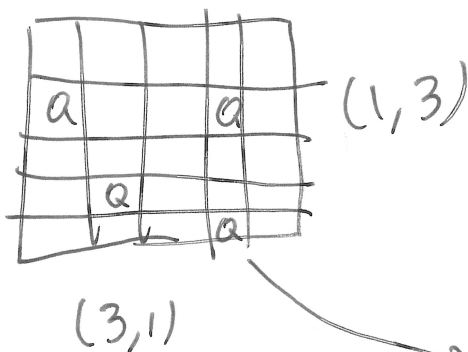
brute force try all perm of $(0,1,2,3)$
place queen on row i , col $\text{perm}[i]$
 $(1,3,0,2)$



$(0,1,-,-,-,-,-)$
 $6!$ wasted work

Idea skip placing 1 in slot 1, since it's doomed to fail.

before placing i in slot k , we check if it conflicts with any previously placed queen.



$$\text{slope} = \frac{3-1}{1-3} = -1 \checkmark$$

$$\text{abs}(\partial x) == \text{abs}(\partial y)$$

$(3,1)$

$(1,0) \quad (4,3) \quad \frac{3-0}{4-1} = 1 \checkmark$

Digit Divisibility

n digit divisible if each of its
 i digit prefixes is ~~div~~ divisible by i
6 digit divisible ~~is~~ is

126402

13 $\rightarrow$ nothing can
start w/13.

1 is divis 1
12 is divis 2
126 is divis 3
1264 is div 4
12640 is div 5
126402 is div 6

