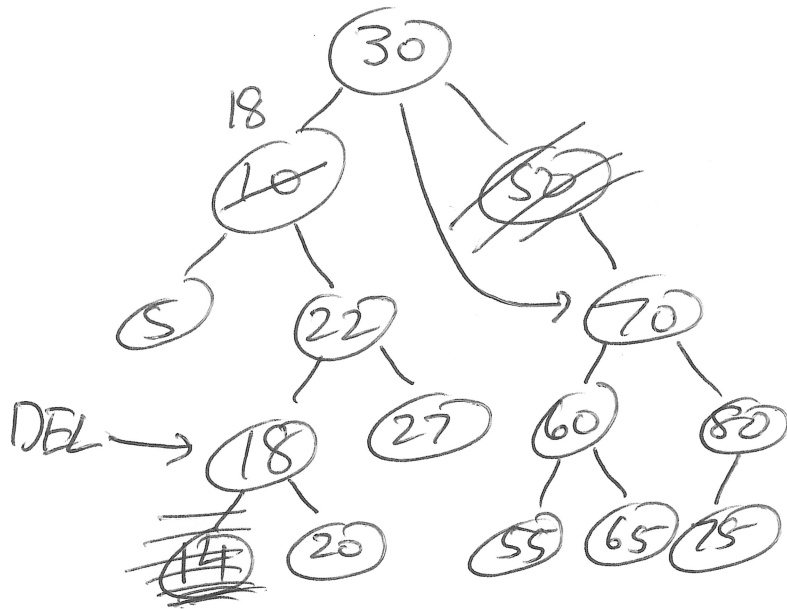


Binary Search Tree Delete BST Problems



Delete 14 (leaf node)

- easy case
- set appropriate ptr (par → left or par → right)
- NULL, free memory
- return value.

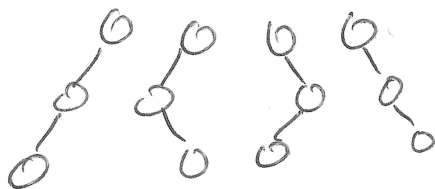
If tree 1 → 0 special case ptr root ⇒ NULL

Delete 50 (one child)

- patch appropriate parent ptr
- appropriate child ptr
- free memory, return value.

- 4 cases

2 child case



Delete 10 (two child)

- Can't patch parent to both children!

Step 1: Copy either max left or min on right into the node storing the value to be deleted.

Step 2: ID node that used to have that value + physically delete it. This node has ≤ 1 child
OLD code can handle it!
min on right has no left child
max of left has no right child.

Storing Extra Data in a Node

default BST:

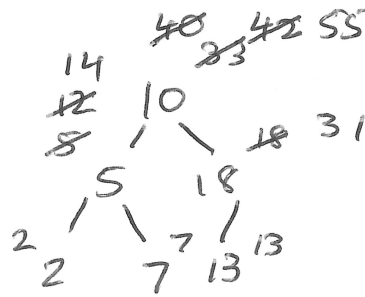
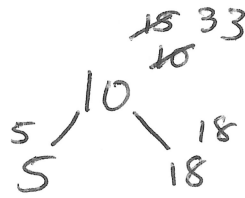
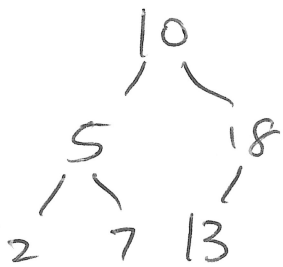
data
left
right

More Stuff

data
sum // sum of all values in subtree
height // the height of this subtree
left
right

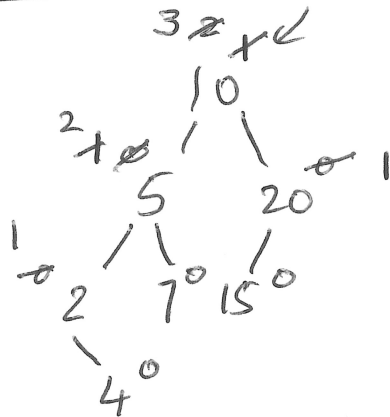
What changes?

In implementation you have to update all the extra data as you insert/delete!

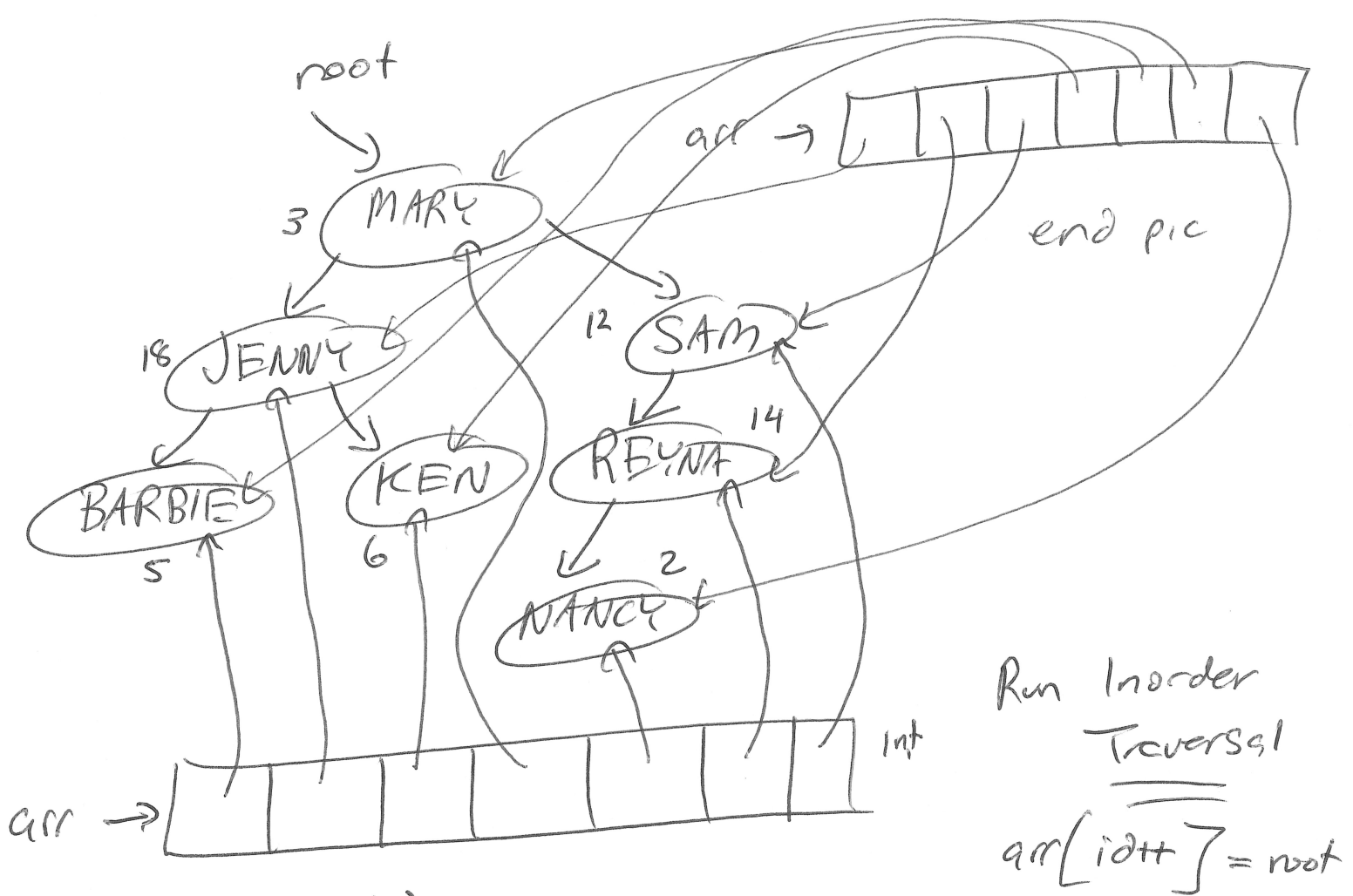


Maintain
Sum

10°



Maintain height



global id
 Sort(arr, n);

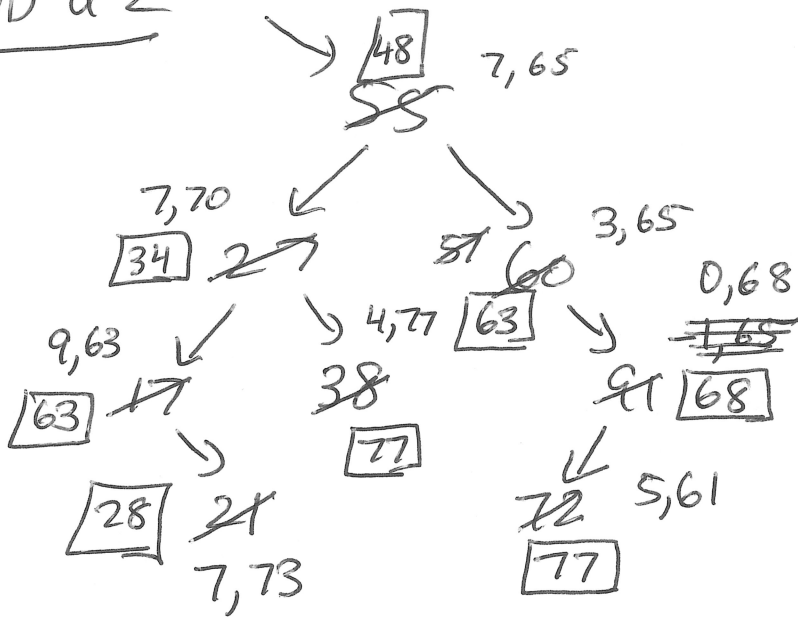
FIND Q1 Notes

$0 \bmod 3$ print data go left
 $1 \bmod 3$ print data + 2 go right
 $2 \bmod 3$ do L, R no print

10 16 30 70 90 85
 ↓ ↓ ↓ ↓ ↓ ↓
12, 18, 30 72, 90, 87

FND Q2

root, 7, 65



FND Q3

```
int check(treenode * A, treenode * B) {
    if (A == NULL && B == NULL) return 1;
    if (A == NULL || B == NULL) return 0;
    if (A->data != B->data) return 0;
    return check(A->left, B->left) && check(A->right, B->right);
}
```