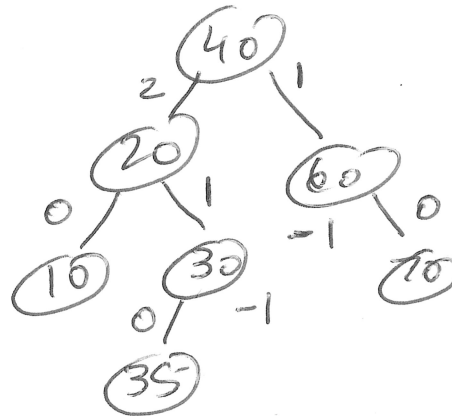


AVL Trees

- Binary Search Tree

- each node follow the avl tree node property
 $abs(\text{height}(\text{left subtree}) - \text{height}(\text{right subtree})) \leq 1$



Any valid AVL tree w/ n nodes has height $O(\lg n)$.

The least number of nodes of an AVL tree of height h is $F_{h+3} - 1$. (F_k is the k^{th} Fib #)

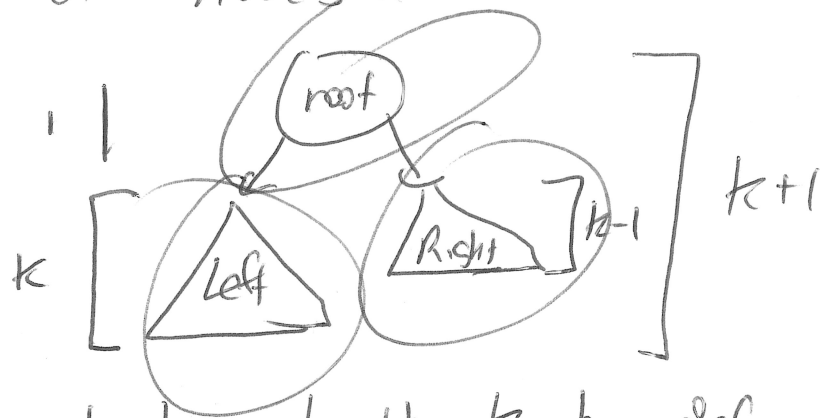
$h=0, \quad \textcircled{20} \rightarrow 1 \text{ node}, \quad F_{0+3} - 1 = 2 - 1 = 1 \quad \checkmark$

$h=1, \quad \begin{matrix} \textcircled{20} \\ \textcircled{30} \end{matrix} \rightarrow 2 \text{ nodes}, \quad F_{1+3} - 1 = 3 - 1 = 2 \quad \checkmark$

1. If assume true for all non-neg int h, $h \leq k$,
 k is a pos int, arbitrarily chosen ≥ 1 . $T(k) = F_{k+3} - 1$.

1.5 Prove for $h = k+1$ that the fewest # nodes in an AVL tree of height $k+1$ is $F_{k+4} - 1$.

Imagine an AVL tree of height $k+1$ w/ the minimal # of nodes:



max subtree must have height k by def.

Other subtree must be at least height $k-1$ by AVL tree property. (Let $T(k) = \min$ # of nodes in AVL tree of height k .)

$$T(k+1) = T(k) + T(k-1) + 1$$

$$= F_{k+3} - 1 + F_{k+2} - 1 + 1$$

$$= F_{k+4} - 1$$

$a=b$ equal

$b=a+b$

$$T(h) = F_{h+3} - 1$$

node $\leftarrow n \geq F_{h+3} - 1$

$$\boxed{\log_{\phi} 2} \approx 1.?$$

Golden ratio

$\phi \leftarrow F_{h+3} \leq n+1$

$$\frac{1}{\sqrt{5}} \left(\left(\frac{1+\sqrt{5}}{2} \right)^{h+3} - \left(\frac{1-\sqrt{5}}{2} \right)^{h+3} \right) \leq n+1$$

$$\frac{1}{\sqrt{5}} \phi^{h+3} \leq n+1$$

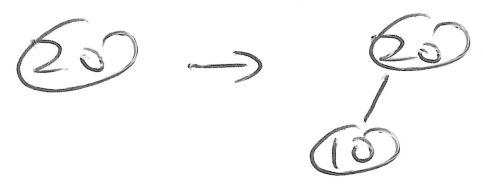
$$\phi^{h+3} \leq c(n+1)$$

$$h+3 \leq \log_{\phi}(c(n+1))$$

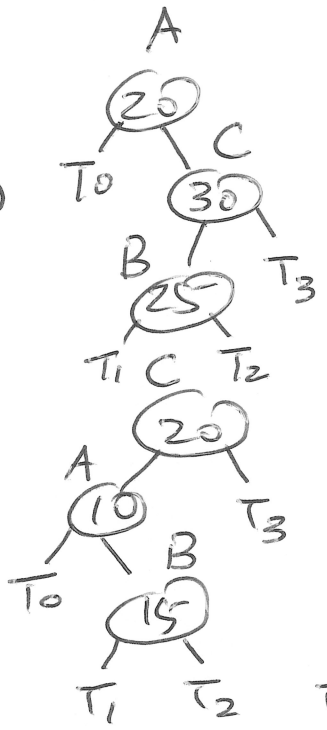
$$\left. \begin{aligned} & \rightarrow h \leq \log_{\phi} c(n+1) - 3 \\ & \leq c \log_{\phi}(n) \end{aligned} \right\}$$

How do we maintain

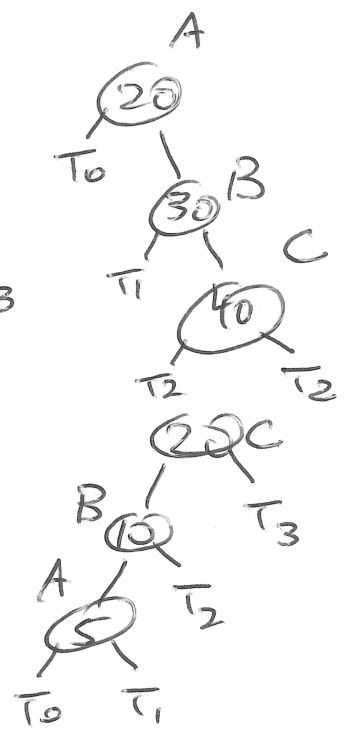
Insert



Double Rotate

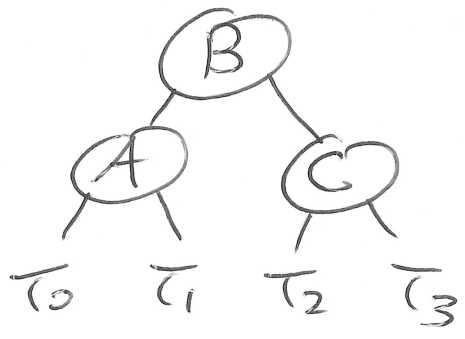


Single Rotate



FIX FOR ALL

4 :

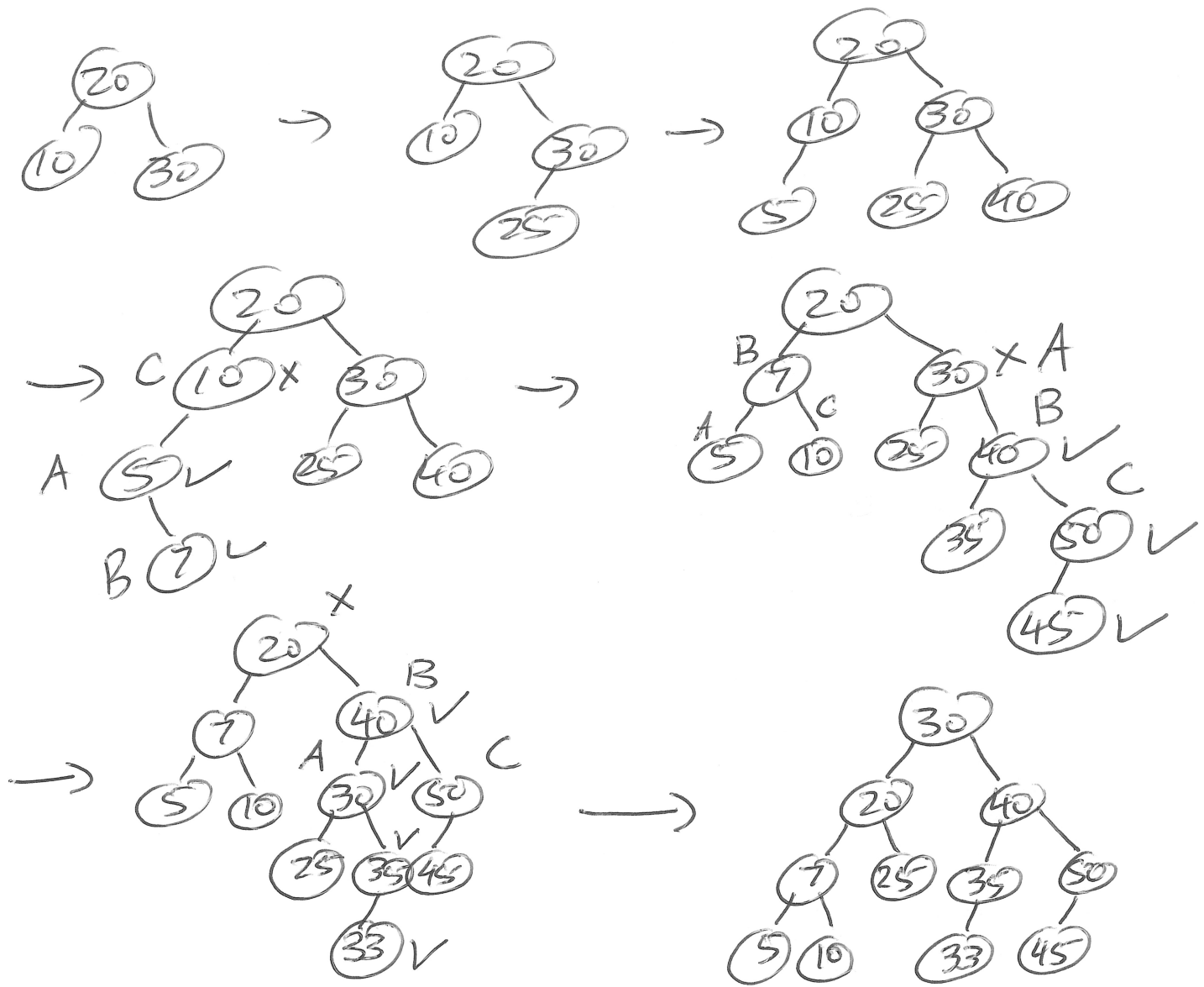


Insert Recursively

root->left = insert(...
look at new height

- a) Do regular bst code
- b) update left + right heights (only 1 of these will be updated)
- c) update my height
- d) If there's an imbalance, fix it.

in code I check this each time, but it never gets called more than once



Delete

Same exact structural code as insert

(a) do regular delete

(b) update height

(c) if imbalance, restructure \Rightarrow might trigger more than

(d) return new root

once!

