

- ① Bubble Sort worst best comparison based Avg
 $O(n^2)$ $O(n^2)$ sorts n items $O(n^2)$
- ② Insertion Sort $O(n^2)$ $O(n)$ $O(n^2)$
- ③ Selection Sort $O(n^2)$ $O(n^2)$ $O(n^2)$
- ④ Merge Sort $O(n \lg n)$ $O(n \lg n)$ $O(n \lg n)$

Bubble Sort

6, 2, 9, 3, 1, 5, 8, 7

2 6 9 3 1 5 8 7

2 6 9 3 1 5 8 7

2 6 3 9 1 5 8 7

2 6 3 1 9 5 8 7

2 6 3 1 5 9 8 7

2 6 3 1 5 8 9 7

2 6 3 1 5 8 7 9

2 3 1 5 6 7 8 9

2 1 3 5 6 7 8 9

1 2 3 5 6 7 8 9

| iteration of
Bubble Sort
after it's done
max element
is in the right
spot.

1st iter

2nd iter

3rd iter

4th iter

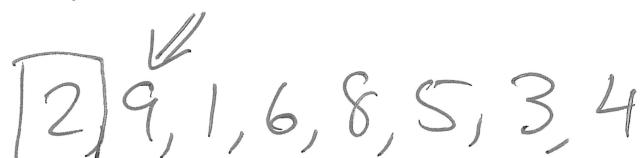
5th, 6th, 7th

Sketch of code
 for (`int j = n-1; j > 0; j--`)

```
for (int i = 0; i < j; i++)
    if (a[i] > a[i+1])
        swap (&a[i], &a[i+1])
```

$$\text{Run time} = (n-1) + (n-2) + (n-3) + \dots + 1 = \frac{(n-1)n}{2} = O(n^2)$$

Insertion Sort



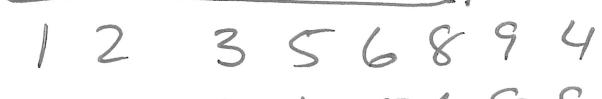














AVG CASE

Iteration i

does $\frac{i}{2}$ swaps

$$\sum_{i=1}^{n-1} \frac{i}{2} = \frac{(n-1)n}{4} = O(n^2)$$

for (`int i = 1; i < n; i++`) {

```
    int j = i - 1;
    while (j > 0 && a[j] < a[j-1]) {
        swap (&a[j], &a[j-1]);
        j--;
    }
}
```

Selection Sort

2, <u>9</u> , 3, 1, 6, 4, 5, 8, <u>7</u>	maxi <u>9</u> 1
2 7 3 1 6 4 5 <u>8</u> 9	1 st iter
2 <u>7</u> 3 1 6 4 5 <u>8</u> 9	2 nd iter (Swapped w/ itself)
2 5 3 1 <u>6</u> 4 7 <u>8</u> 9	3 rd iter
2 <u>5</u> 3 1 4 6 7 <u>8</u> 9	4 th iter
2 <u>4</u> 3 1 5 6 7 <u>8</u> 9	5 th iter
2 1 <u>3</u> 4 5 6 7 <u>8</u> 9	6 th iter
<u>2</u> 1 3 4 5 6 7 <u>8</u> 9	7 th iter
1 2 3 4 5 6 7 8 9	8 th iter

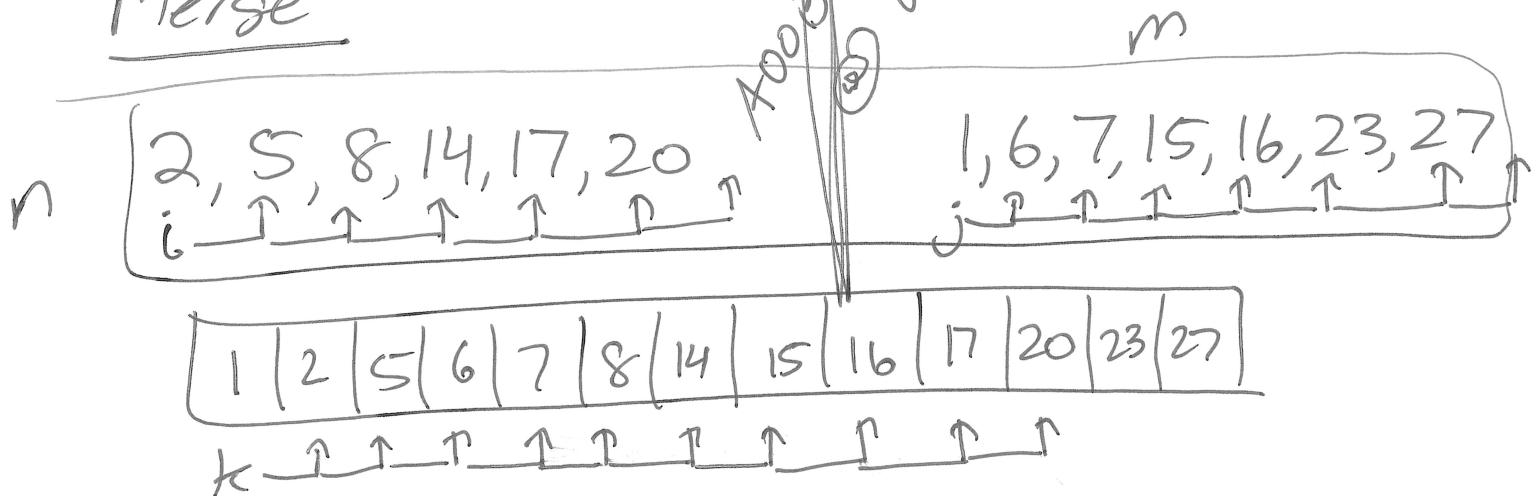
```

for (int j=n-1; j>0; j--) {
    int maxi = 0;
    for (int i = 0; i<=j; i++)
        if ( a[i] > a[maxi] )
            maxi = i;
    swap( &a[maxi], &a[j] );
}

```

Merge Sort

Merge



Run-Time : $O(n+m)$, $O(\max(n,m))$

MergeSort (int arr, int SI, int EI) {

if ($SI \geq EI$) return;

int mid = $(SI + EI) / 2$;

→ MergeSort (arr, SI, mid);

MergeSort (arr, mid+1, EI);

Merge (arr, SI, mid+1, EI); // Pass in 2 arrays

}

Sort Left Sort Right → merge

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + O(n)$$

$$= 2T\left(\frac{n}{2}\right) + O(n)$$

Via Iteration, Master Thm $T(n) = O(n \lg n)$