

Simplifying Reasoning about Objects with Tako

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Immediate objective

Introduce students to “heavyweight” formal reasoning without scaring them (too much)

- familiar syntax (Java-like)
- value semantics
- variable-based state space

The Tako Language

- Similar to Java, but...
 - **Avoids common sources of aliasing**
 - Automatically initializes variables
 - Uses swapping ($:=:$) instead of assignment
 - Methods have in-out parameter passing


Tako

```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
        MAX := n;
        contents := new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] ::= x;
        top++;
    }
    public void pop(Object x) {
        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
    }
}
```

Java

```
public class BddStack {
    private final int MAX;
    private Object[ ] contents;
    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```

Tako



```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
        MAX := n;
        contents := new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] ::= x;
        top++;
    }
    public void pop(Object x) {
        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
    }
}
```

Java

```
public class BddStack {
    private final int MAX;
    private Object[ ] contents;
    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```

Tako Initialization

- By convention, Tako programmers should always write a default constructor, which is applied to all variables as soon as they are declared.
- `Integer top; ≡ Integer top := new Integer();`
- `IStack s; ≡ IStack s := null; // interface`
- But use `BddStack s := new BddStack(20);`


Tako

```
public class BddStack {  
    private final Integer MAX;  
    private Object[ ] contents;  
    private Integer top;  
    public BddStack(Integer n) {  
        MAX := n;  
        contents := new Object[MAX];  
    }  
    public void push(Object x) {  
        assert depth( ) < MAX;  
        contents[top] ::= x;  
        top++;  
    }  
    public void pop(Object x) {  
        assert depth( ) > 0;  
        top--;  
        x ::= contents[top];  
    }  
    public Integer depth( ) {  
        result := top;  
    }  
}
```

Java

```
public class BddStack {  
    private final int MAX;  
    private Object[ ] contents;  
    private int top;  
    public BddStack(int n) {  
        MAX = n;  
        contents = new Object[MAX];  
    }  
    public void push(Object x) {  
        assert depth( ) < MAX;  
        contents[top] = x;  
        top++;  
    }  
    public Object pop( ) {  
        assert depth( ) > 0;  
        top--;  
        return contents[top];  
    }  
    public int depth( ) {  
        return top;  
    }  
}
```

Tako



```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
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        contents := new Object[MAX];
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    public void push(Object x) {
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        contents[top] ::= x;
        top++;
    }
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        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
    }
}
```

Java

```
public class BddStack {
    private final int MAX;
    private Object[ ] contents;
    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```


Function Assignment

- Has the form `variable := my.expression();`
- If the compiler finds:
`MAX := n;`
it will translate it as:
`MAX := n.replica();`
where `replica()` is a programmer defined method expected to perform a deep copy

Tako

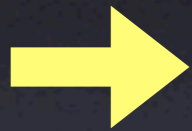
```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
        MAX := n;
        contents := new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] ::= x;
        top++;
    }
    public void pop(Object x) {
        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
    }
}
```

Java

```
public class BddStack {
    private final int MAX;
    private Object[ ] contents;
    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
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    public int depth( ) {
        return top;
    }
}
```

Tako

```
public class BddStack {
    private final Integer MAX;
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        assert depth( ) > 0;
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        result := top;
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Java

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public class BddStack {
    private final int MAX;
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    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
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    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```

Swapping

- Compiler can **implement** it by **swapping references**
- Preserves unique references, so the programmer can **reason** about it as if **objects are swapped**

Tako


```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
        MAX := n;
        contents := new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] ::= x;
        top++;
    }
    public void pop(Object x) {
        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
    }
}
```

Java

```
public class BddStack {
    private final int MAX;
    private Object[ ] contents;
    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
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Tako

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public class BddStack {
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        contents[top] ::= x;
        top++;
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    public void pop(Object x) {
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Java

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public class BddStack {
    private final int MAX;
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    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
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    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```

Tako Parameter Passing

- Unlike Java, variable values in Tako can always be updated through parameters
- Similar to:
 - C++ pass-by-reference
 - C# ref parameters
 - Ada in-out parameter passing
 - Pascal value-result parameter passing

Functions vs. Procedures

- By convention, Tako functions (non-void methods) do not change the program state
(for all variables v , $v = \#v$)
- If you want to change a variable's value, use a procedure (void method)
- For example, use **public void** pop(Item x); instead of **public** Item pop();

Tako


```
public class BddStack {
    private final Integer MAX;
    private Object[ ] contents;
    private Integer top;
    public BddStack(Integer n) {
        MAX := n;
        contents := new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] ::= x;
        top++;
    }
    public void pop(Object x) {
        assert depth( ) > 0;
        top--;
        x ::= contents[top];
    }
    public Integer depth( ) {
        result := top;
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}
```

Java

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public class BddStack {
    private final int MAX;
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    private int top;
    public BddStack(int n) {
        MAX = n;
        contents = new Object[MAX];
    }
    public void push(Object x) {
        assert depth( ) < MAX;
        contents[top] = x;
        top++;
    }
    public Object pop( ) {
        assert depth( ) > 0;
        top--;
        return contents[top];
    }
    public int depth( ) {
        return top;
    }
}
```

Tako

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public class BddStack {  
    private final Integer MAX;  
    private Object[ ] contents;  
    private Integer top;  
    public BddStack(Integer n) {  
        MAX := n;  
        contents := new Object[MAX];  
    }  
    public void push(Object x) {  
        assert depth( ) < MAX;  
        contents[top] ::= x;  
        top++;  
    }  
    public void pop(Object x) {  
        assert depth( ) > 0;  
        top--;  
        x ::= contents[top];  
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    public Integer depth( ) {  
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Java

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public class BddStack {  
    private final int MAX;  
    private Object[ ] contents;  
    private int top;  
    public BddStack(int n) {  
        MAX = n;  
        contents = new Object[MAX];  
    }  
    public void push(Object x) {  
        assert depth( ) < MAX;  
        contents[top] = x;  
        top++;  
    }  
    public Object pop( ) {  
        assert depth( ) > 0;  
        top--;  
        return contents[top];  
    }  
    public int depth( ) {  
        return top;  
    }  
}
```

```
public Integer depth( ) {  
    result := top;  
}
```

translates to

```
public Integer depth( ) {  
    Integer result = new Integer( );  
    result = top.replica( );  
    return result;  
}
```

Clean semantics

Variable-based
property

The **state space** is simply the abstract **object values** of the defined variables

Frame
property

During a **method call**, the only **variable values that may change** are:

1. arguments to the call, and
2. globals defined in the *affects* clause.

```
import spec.MathString;

public interface Stack {
    model MathString;

    initialization ensures
        this = EMPTY_STRING;

    public void push(Object x);
        ensures this = <#x> 0 #this;

    public void pop(Object x);
        requires |this| > 0;
        ensures #this = <x> 0 this;

    public Integer depth( );
        ensures this = #this and result = |this|;
}
```

```
import spec.MathString;

public interface Stack {
    model MathString;

    initialization ensures
        this = EMPTY_STRING;

    public void push(Object x);
        ensures this = <#x> 0 #this;

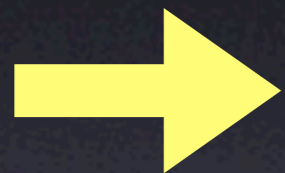
    public void pop(Object x);
        requires |this| > 0;
        ensures #this = <x> 0 this;

    public Integer depth( );
        ensures this = #this and result = |this|;
}
```

```
import spec.MathString;
```

```
public interface Stack {  
    model MathString;
```

```
    initialization ensures  
        this = EMPTY_STRING;
```



```
    public void push(Object x);  
        ensures this = <#x> 0 #this;
```

```
    public void pop(Object x);  
        requires |this| > 0;  
        ensures #this = <x> 0 this;
```

```
    public Integer depth( );  
        ensures this = #this and result = |this|;
```

```
}
```

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> 0 #**this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> 0 #this;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this ≡ <#x> 0 #**this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> 0 **#this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> o #**this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> 0 **#this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value

ensures this = <#x> 0 #**this**;

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value and the current x value is a valid but unspecified value

ensures this = <#x> 0 #**this**;

and x = ??

public void push(Object x) guarantees that . . .

The current stack value is equal to a string containing the old x value concatenated with the old stack value and the current x value is a valid but unspecified value and no other variables in the state space change

ensures this = <#x> 0 #**this**;

and x = ??

and v1 = #v1 **and** v2 = #v2 **and** . . .

Three most important facts for Java programmers to remember about Tako

1. Everything is *an object*
2. *Everything* is an object
3. See facts 1 & 2

Reasoning about a Tako Stack

$s = \langle \text{[graph1]}, \text{[graph2]}, \text{[graph3]} \rangle$ and $g = \text{[graph1]}$

`g.deleteANode();`

$s = \langle \text{[graph1]}, \text{[graph2]}, \text{[graph3]} \rangle$ and $g = \text{[graph4]}$

Reasoning about a Java Stack

$s = \langle \text{[graph]}, \text{[edge]}, \text{[triangle]} \rangle$ and $g = \text{[graph]}$

`g.deleteANode();`

$s = \langle \text{[?]}, \text{[edge]}, \text{[triangle]} \rangle$ and $g = \text{[graph]}$

$\text{[?]} = \text{[graph]} \text{ or } \text{[graph]}$

Pointer Component

- A pointer component is provided for times when programmers need the behavior and performance of pointer (as in implementing *linked data structures*)
- It is similar to the RESOLVE pointer component (see Kulczycki in SAVCBS2005)
- Model uses *conceptual variables*

Tracing Table for a Method

- **Given:** an example program state that conforms to the method's precondition
- **After each statement:** update the state based on the statement's specification
- **Final state:** should be consistent with the specification of the original method

```
public void reverse( )  
    ensures this = REV(#this);  
{  
    Stack temp;  
    Object x;  
    while (this.depth( ) != 0)  
        decreasing . . .  
        maintaining . . .  
    {  
        this.pop(x);  
        temp.push(x);  
    }  
    this ::= temp;  
}
```

pre-state: **this** = < 3, 4 >



```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

expected post-state: **this** = < 3, 4 >

pre-state: **this** = < 3, 4 >

0 **this** = < 3, 4 > **and** temp = < > **and** x = 0

1

2

3

4

5

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0 **this** = < 3, 4 > **and** temp = < > **and** x = 0

1 **this** = < 3, 4 > **and**
temp = < > **and**
x = 0

2

3

4

5

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0 **this** = < 3, 4 > **and** temp = < > **and** x = 0

1 **this** = < 3, 4 > **and**
temp = < > **and**
x = 0

2 **this** = < 4 > **and**
temp = < > **and**
x = 3

3

4

5

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0 **this** = < 3, 4 > **and** temp = < > **and** x = 0

1 **this** = < 3, 4 > **and**
temp = < > **and**
x = 0

2 **this** = < 4 > **and**
temp = < > **and**
x = 3

3 **this** = < 4 > **and**
temp = < 3 > **and**
x = ??

4

5

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0	this = < 3, 4 > and temp = < > and x = 0
1	this = < 3, 4 > and temp = < > and x = 0
2	this = < 4 > and temp = < > and x = 3
3	this = < 4 > and temp = < 3 > and x = ??
4	
5	

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
} // this.depth() != 0
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0	this = < 3, 4 > and temp = < > and x = 0	
1	this = < 3, 4 > and temp = < > and x = 0	this = < 4 > and temp = < 3 > and x = ??
2	this = < 4 > and temp = < > and x = 3	this = < > and temp = < 3 > and x = 4
3	this = < 4 > and temp = < 3 > and x = ??	this = < > and temp = < 4, 3 > and x = ??
4		
5		

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

expected post-state: **this** = < 3, 4 >

pre-state: **this** = < 3, 4 >

0	this = < 3, 4 > and temp = < > and x = 0	
1	this = < 3, 4 > and temp = < > and x = 0	this = < 4 > and temp = < 3 > and x = ??
2	this = < 4 > and temp = < > and x = 3	this = < > and temp = < 3 > and x = 4
3	this = < 4 > and temp = < 3 > and x = ??	this = < > and temp = < 4, 3 > and x = ??
4		
5		

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
} // this.depth() = 0
```

```
this ::= temp;
```

pre-state: **this** = < 3, 4 >

0	this = < 3, 4 > and temp = < > and x = 0	
1	this = < 3, 4 > and temp = < > and x = 0	this = < 4 > and temp = < 3 > and x = ??
2	this = < 4 > and temp = < > and x = 3	this = < > and temp = < 3 > and x = 4
3	this = < 4 > and temp = < 3 > and x = ??	this = < > and temp = < 4, 3 > and x = ??
4	this = < > and temp = < 4, 3 > and x = ??	
5		

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```


pre-state: **this** = < 3, 4 >

0	this = < 3, 4 > and temp = < > and x = 0	
1	this = < 3, 4 > and temp = < > and x = 0	this = < 4 > and temp = < 3 > and x = ??
2	this = < 4 > and temp = < > and x = 3	this = < > and temp = < 3 > and x = 4
3	this = < 4 > and temp = < 3 > and x = ??	this = < > and temp = < 4, 3 > and x = ??
4	this = < > and temp = < 4, 3 > and x = ??	
5	this = < 4, 3 > and temp = < > and x = ??	

expected post-state: **this** = < 3, 4 >

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

Symbolic Reasoning Table

- Used to prove an implementation is correct with respect to its specification
- Assume precondition; show postcondition
- For each state you have:
 - **Path condition** – What has to be true in order for us to get to this state?
 - **Facts** – What do we know in this state?
 - **Obligations** – What do we need to prove before we can move to the next state?

Proof of Correctness

- To prove an implementation correct with respect to its specification:
 - Construct a symbolic reasoning table
 - Prove that all the obligations are true

```
public void reverse( )  
    ensures this = REV(#this);  
{  
    Stack temp;  
    Object x;  
    while (this.depth( ) != 0)  
        decreasing |this|;  
        maintaining REV(temp) o this = #this;  
    {  
        this.pop(x);  
        temp.push(x);  
    }  
    this ::= temp;  
}
```

PC

Facts

Obligations

0

1

2

3

4

5

```
while (this.depth() != 0) {
```

```
this.pop(x);
```

```
temp.push(x);
```

```
}
```

```
this ::= temp;
```

PC

Facts

Obligations

0

1

2

3

4

5

maintaining
REV(temp) o **this**
= #**this**

```
while (this.depth() != 0) {
```

```
this.pop(x);
```

```
temp.push(x);
```

```
}
```

```
this ::= temp;
```

PC Facts Obligations

0		
1	REV(temp ₁) o this = # this	
2		
3		REV(temp ₃) o this = # this
4		
5		

maintaining
REV(temp) o **this**
= #**this**

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC Facts Obligations

0		
1	REV(temp ₁) o this ₁ = # this	
2		
3		REV(temp ₃) o this ₃ = # this
4		
5		

decreasing |this|

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC

Facts

Obligations

decreasing |this|

0

1

2

3

4

5

REV(temp₁) o **this**₁
= #**this** and
x₁ = ??

REV(temp₃) o **this**₃
= #**this** and
|**this**₃| < |**this**₁|

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC

Facts

Obligations

0

1

2

3

4

5

REV(temp₁) o **this**₁
= #**this** and
x₁ = ??

REV(temp₃) o **this**₃
= #**this** and
|**this**₃| < |**this**₁|

|**this**_{pre_1}| = 0 and
this₁ = **this**_{pre_1}

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC

Facts

Obligations

0

1

2

3

4

5

$|this_{pre_1}| = 0$ and $this_1 = this_{pre_1}$

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

$ this_1 \neq 0$	REV(temp ₁) o this ₁ = # this and x ₁ = ??	
$ this_1 \neq 0$		
$ this_1 \neq 0$		REV(temp ₃) o this ₃ = # this and $ this_3 < this_1 $

PC

Facts

Obligations

requires
|this| > 0

0

1

2

3

4

5

 this₁ ≠ 0	REV(temp ₁) o this₁ = #this and x ₁ = ??	
 this₁ ≠ 0		
 this₁ ≠ 0		REV(temp ₃) o this₃ = #this and this₃ < this₁

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC

Facts

Obligations

requires
|this| > 0

0

1

2

3

4

5

 this₁ ≠ 0	REV(temp ₁) o this₁ = #this and x ₁ = ??	 this₁ > 0
 this₁ ≠ 0		
 this₁ ≠ 0		REV(temp ₃) o this₃ = #this and this₃ < this₁

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC

Facts

Obligations

0

1

2

3

4

5

**ensures #this =
<x> 0 this**

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

$ this_1 \neq 0$	REV(temp ₁) o this ₁ = # this and x ₁ = ??	$ this_1 > 0$
$ this_1 \neq 0$		
$ this_1 \neq 0$		REV(temp ₃) o this ₃ = # this and $ this_3 < this_1 $

PC

Facts

Obligations

0		
1	$ \mathbf{this}_1 \neq 0$ REV(temp ₁) o this ₁ = # this and x ₁ = ??	$ \mathbf{this}_1 > 0$
2	$ \mathbf{this}_1 \neq 0$ this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
3	$ \mathbf{this}_1 \neq 0$	REV(temp ₃) o this ₃ = # this and $ \mathbf{this}_3 < \mathbf{this}_1 $
4		
5		

ensures #this =
<x> o this

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

-- no requires clause for push

	PC	Facts	Obligations
0			
1	$ \mathbf{this}_1 \neq 0$	REV(temp ₁) o this ₁ = # this and x ₁ = ??	$ \mathbf{this}_1 > 0$
2	$ \mathbf{this}_1 \neq 0$	this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
3	$ \mathbf{this}_1 \neq 0$		REV(temp ₃) o this ₃ = # this and $ \mathbf{this}_3 < \mathbf{this}_1 $
4			
5			

```

while (this.depth() != 0) {
    this.pop(x);
    temp.push(x);
}
this ::= temp;

```


PC

Facts

Obligations

0

1

2

3

4

5

$ \mathbf{this}_1 \neq 0$	REV(temp ₁) o this ₁ = # this and x ₁ = ??	$ \mathbf{this}_1 > 0$
$ \mathbf{this}_1 \neq 0$	this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
$ \mathbf{this}_1 \neq 0$		REV(temp ₃) o this ₃ = # this and $ \mathbf{this}_3 < \mathbf{this}_1 $

ensures this =
<#x> o #this

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC**Facts****Obligations**

0		
1	$ \mathbf{this}_1 \neq 0$ $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \#\mathbf{this}$ and $x_1 = ??$	$ \mathbf{this}_1 > 0$
2	$ \mathbf{this}_1 \neq 0$ $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$	
3	$ \mathbf{this}_1 \neq 0$ $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$	$\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \#\mathbf{this}$ and $ \mathbf{this}_3 < \mathbf{this}_1 $
4		
5		

ensures this =
 $\langle \#x \rangle \circ \#\mathbf{this}$

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

-- other parts deal with initialization and termination of loop invariant

	PC	Facts	Obligations
0			
1	$ \mathbf{this}_1 \neq 0$	REV(temp ₁) o this ₁ = # this and x ₁ = ??	$ \mathbf{this}_1 > 0$
2	$ \mathbf{this}_1 \neq 0$	this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
3	$ \mathbf{this}_1 \neq 0$	this ₃ = <x ₂ > o this ₂ and x ₃ = ?? and temp ₃ = temp ₂	REV(temp ₃) o this ₃ = # this and $ \mathbf{this}_3 < \mathbf{this}_1 $
4			
5			

while (**this**.depth() != 0) {

this.pop(x);

temp.push(x);

}

this ::= temp;

PC**Facts****Obligations**

0	Object.is_init(x ₀) and temp ₀ = < >	this ₀ ≠ 0 implies REV(temp ₀) o this ₀ = # this
1	this ₁ ≠ 0 REV(temp ₁) o this ₁ = # this and x ₁ = ??	this ₁ > 0
2	this ₁ ≠ 0 this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
3	this ₁ ≠ 0 this ₃ = <x ₂ > o this ₂ and x ₃ = ?? and temp ₃ = temp ₂	REV(temp ₃) o this ₃ = # this and this ₃ < this ₁
4	this ₄ = 0 REV(temp ₄) o this ₄ = # this and x ₄ = ??	
5	this ₄ = 0 this ₅ = temp ₄ and temp ₅ = this ₄ and x ₅ = x ₄	this ₅ = REV(# this)

```
while (this.depth() != 0) {
```

```
  this.pop(x);
```

```
  temp.push(x);
```

```
}
```

```
this ::= temp;
```

PC Facts Obligations

0	Object.is_init(x ₀) and temp ₀ = < >	this ₀ ≠ 0 implies REV(temp ₀) o this ₀ = #this
1	this ₁ ≠ 0 REV(temp ₁) o this ₁ = #this and x ₁ = ??	this ₁ > 0
2	this ₁ ≠ 0 this ₁ = <x ₂ > o this ₂ and temp ₂ = temp ₁	
3	this ₁ ≠ 0 this ₃ = <x ₂ > o this ₂ and x ₃ = ?? and temp ₃ = temp ₂	REV(temp ₃) o this ₃ = #this and this ₃ < this ₁
4	this ₄ = 0 REV(temp ₄) o this ₄ = #this and x ₄ = ??	
5	this ₄ = 0 this ₅ = temp ₄ and temp ₅ = this ₄ and x ₅ = x ₄	this ₅ = REV(#this)

while (this.depth() != 0) {

 this.pop(x);

 temp.push(x);

}

this ::= temp;

Show: $\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \#\mathbf{this}$

Fact 1: $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \#\mathbf{this}$ and $x_1 = ??$

Fact 2: $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$

Fact 3: $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$

Show: $\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \#\mathbf{this}$

or

$\text{REV}(\text{temp}_2) \circ \langle x_2 \rangle \circ \mathbf{this}_2 = \#\mathbf{this}$ (by Fact 3)

Fact 1: $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \#\mathbf{this}$ and $x_1 = ??$

Fact 2: $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$

Fact 3: $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$

Show: $\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \#\mathbf{this}$

or

$\text{REV}(\text{temp}_2) \circ \langle x_2 \rangle \circ \mathbf{this}_2 = \#\mathbf{this}$ (by Fact 3)

or

$\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \#\mathbf{this}$ (by Fact 2)

Fact 1: $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \#\mathbf{this}$ and $x_1 = ??$

Fact 2: $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$

Fact 3: $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$

Show: $\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \# \mathbf{this}$

or

$\text{REV}(\text{temp}_2) \circ \langle x_2 \rangle \circ \mathbf{this}_2 = \# \mathbf{this}$ (by Fact 3)

or

$\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \# \mathbf{this}$ (by Fact 2)

or

$\# \mathbf{this} = \# \mathbf{this}$ (by Fact 1)

Fact 1: $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \# \mathbf{this}$ and $x_1 = ??$

Fact 2: $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$

Fact 3: $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$

Show: $\text{REV}(\text{temp}_3) \circ \mathbf{this}_3 = \# \mathbf{this}$

or

$\text{REV}(\text{temp}_2) \circ \langle x_2 \rangle \circ \mathbf{this}_2 = \# \mathbf{this}$ (by Fact 3)

or

$\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \# \mathbf{this}$ (by Fact 2)

or

 $\# \mathbf{this} = \# \mathbf{this}$ (by Fact 1)

Fact 1: $\text{REV}(\text{temp}_1) \circ \mathbf{this}_1 = \# \mathbf{this}$ and $x_1 = ??$

Fact 2: $\mathbf{this}_1 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $\text{temp}_2 = \text{temp}_1$

Fact 3: $\mathbf{this}_3 = \langle x_2 \rangle \circ \mathbf{this}_2$ and $x_3 = ??$ and $\text{temp}_3 = \text{temp}_2$

Tako compiler

- **SourceForge** project: **takocompiler**
- Text-based adventure program:
 - ~40 classes
 - ~4000 lines of code

Questions

Inheritance and in-out parameter passing

- ```
void swapMethod(Object x, Object y) {
 x := y;
}
```
- ```
Cat c; Dog d;  
swapMethod(c, d);
```